



UNIVERSITY OF CAPE TOWN  
DEPARTMENT OF COMPUTER SCIENCE

THE IMPLEMENTATION OF A FRONT END PROCESSOR  
FOR A SUBSET OF ADA (1)

BY

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A THESIS

PREPARED UNDER THE SUPERVISION OF  
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IN FULFILMENT OF THE REQUIREMENTS FOR THE  
DEGREE OF MASTER OF SCIENCE IN COMPUTER SCIENCE

(1) ADA IS A REGISTERED TRADEMARK OF THE U.S. GOVERNMENT,  
ADA JOINT PROGRAM OFFICE

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ABSTRACT.

ADA is a high level programming language sponsored by the United States Department of Defence primarily for use in real-time systems. It has all the structures present in modern algorithmic languages with additional features for tasking.

This thesis discusses the University of Cape Town implementation of a front end processor for a subset of ADA. A compiler generator package was used to construct a syntax checker for the ADA language and a subset of this was extended through the semantic analysis phase finally to produce the intermediate code - DIANA.

DIANA is the standard intermediate code for all ADA programs and a representation for transfer between systems has been defined. DIANA is intended to function as an interface between the front and back ends of ADA compilers, and as an intermediate form which can be used by tools designed for ADA.

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## 1. CHAPTER 1. INTRODUCTION.

ADA is a high level programming language initiated by the United States Department of Defence. During its development considerable effort went into the determination of the requirements of a language intended primarily for embedded computer applications. The resulting language is suitable not only for embedded computer applications, but also for general systems applications, numeric computation, real-time industrial applications, general applications programming, and for teaching good programming practices.

The first specification of the language became available in 1979 and in 1980 students at the University of Cape Town decided to gain experience in this new language by attempting to develop an ADA front end processor.

In producing the front end processor, a compiler generator package, GSA 1100, was used. This package provides the capability of writing a lexical analyser (accepting as input, a modified Backus Naur grammar), and an LL(K) parser. It also enables error recovery and calls to semantic routines to be embedded in the parser specification. A description of the GSA 1100 package is contained in CHAPTER 2.

The version of ADA used by the initial project was that released in 1979. This was modified into ADA-80 which became available in 1980, and is the version used by the final project. The 1980 definition did not entirely resolve the ambiguities and complexities of the 1979 specification, and ADA-80 had to be further refined in 1982.

The intermediate code produced by the front end was DIANA, an abstract syntax tree. It was designed for ADA-80 and has been

updated to be compatible with ADA-82. CHAPTER 3 discusses the history and features of these two languages.

While implementing ADA using GSA 1100, certain problems were encountered. Most of these occurred because the ADA specification is ambiguous and contains many constructs unsuited to an LL(K) parser. The ADA specification was therefore restructured to permit parsing using GSA 1100. CHAPTER 4 describes how the syntax checker was constructed, highlighting the problem areas encountered and their solution.

ADA has rigid typing rules and the language specification requires that the static semantic checking be done at compile time. This entails type checking, evaluation of numeric expressions, checking for non-null ranges etc. To achieve this the program symbol table must be referenced. CHAPTER 5 describes the construction of the symbol table, and the routines that reference it while performing the semantic checking.

The output of the front end processor system is an abstract syntax tree representation of the intermediate code DIANA. DIANA was produced jointly by designers at Carnegie-Mellon University and the University of Karlsruhe for the ADA language. A description of DIANA is contained in CHAPTER 6.

CHAPTER 7 contains examples run on the front end processor and the resulting error messages and code produced.

## 2. CHAPTER 2. STRUCTURE OF THE ADA COMPILER.

The University of Cape Town ADA front end processor is written for the UNIVAC 1100/81. The body of the project is written in PLUS, Programming Language for Univac Systems, [Sperry Univac 1] and interfaces with tables generated by a syntax analyser package, GSA 1100, [Sperry Univac 2] provided by UNIVAC. Driver routines and interfaces are written in Univac assembler.

Figure 2-1 shows the basic structure of the envisaged compiler.

The structure of the UCT-front end processor is divided into five phases:

Phases 1, 2 and 3 are accomplished with the aid of the GSA 1100 package. Using the output produced by this package a symbol table is built which stores all tokens and their attributes as they are defined and later modified. This information is then used in the production of DIANA and enables the semantic information to be included in the DIANA structure.

Phase 4, the static semantic analysis involves checking that the tokens are correctly used, and so uses the symbol table produced by phases 1 and 2. If necessary, additional information may be inserted in the symbol table. In the case of errors, the message table produced in phase 3 is used.

In the production of DIANA, phase 5, all the semantic information comes from the symbol table, while the knowledge of which node to generate, or to where a particular node attribute should be linked, comes from the parse table generated in phase 2.

In this way, all the phases making up the front end are interdependent, and use the information produced by each other.

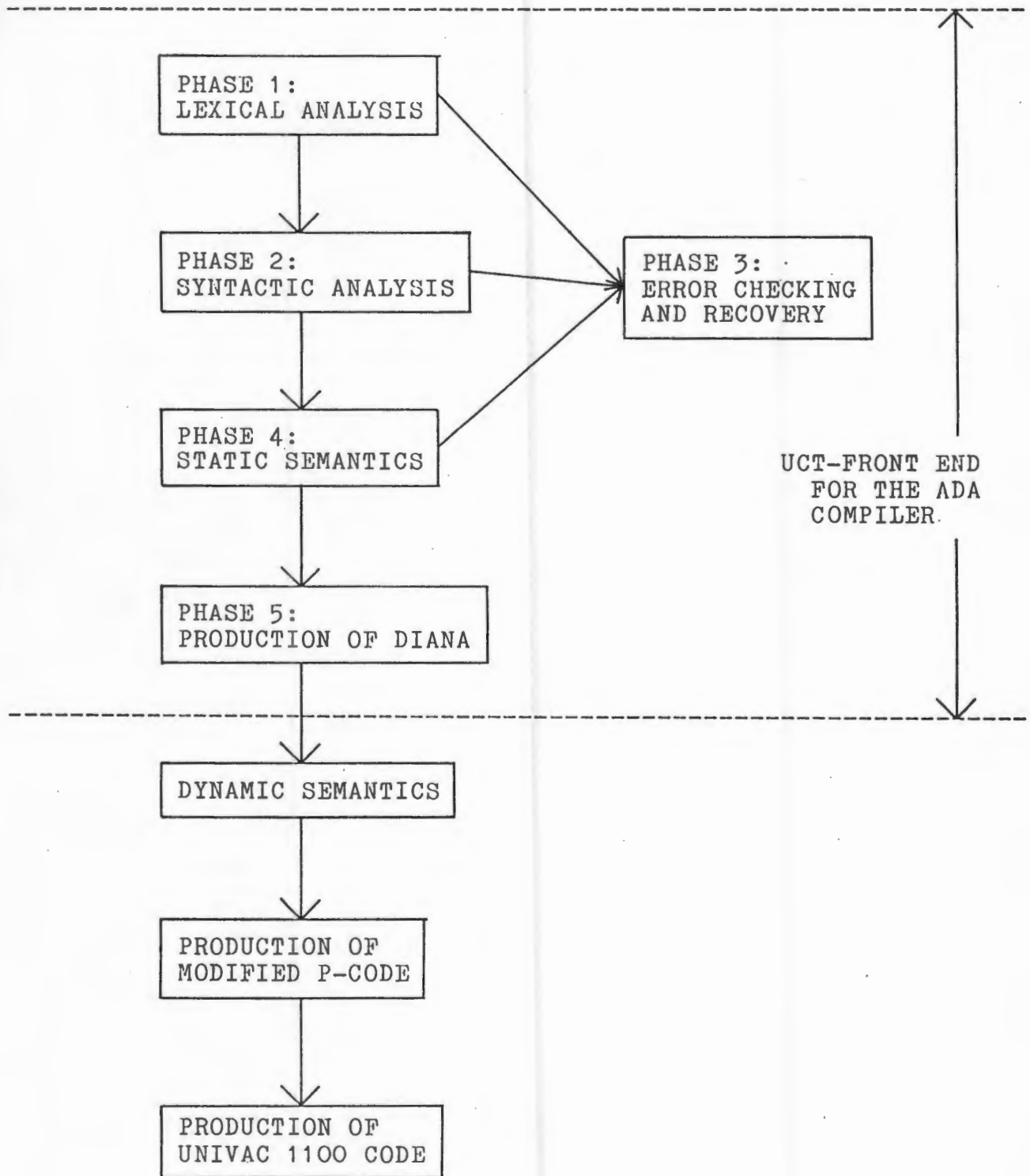


Figure 2-1. STRUCTURE OF THE UCT ADA COMPILER.

## 2.1. TABLE GENERATING ROUTINES.

The general syntax analyser consists of several table generating routines that take as input information about the language to be analysed, and produce as output various tables that may be interfaced with user-written routines. The table generating routines comprise the lexical analyser, parser, error message table and dictionary. Interfacing with each of the above is assembler code which manipulates the output produced by the package and passes tokens to user-written procedures via predefined stacks.

### 2.1.1. LEXICAL ANALYSER.

Input to the lexical analyser is in the form of a modified Backus Naur notation describing the composition of the terminal symbols that make up the lexical units of the language being defined. (See appendix I).

Once a lexical unit is recognised, it is transferred to the parser for parsing.

Each lexical unit is given a unique parameter number in the lexical specification, and this coincides with the same parameter number for the same token in the parser specification.

The lexical specification can be considered as a tree with many branches. It consists of only one production and so has one entry point, the root. Depending on the next character it follows a particular branch until a token is recognised, or an error found. Look-ahead should only be used in the lexical specification when absolutely necessary. This implementation had to use look-ahead to distinguish between a character literal e.g. 'R' and an

attribute reference e.g. A'FIRST. After recognising a single quote mark followed by a character, the path to follow depends on the next character in the input string.

### 2.1.2. PARSER.

Unlike the lexical analyser, the parser has many productions which make up the specification of the language. The same notation is used to input the rules of the language to the parser as is used to input the lexical units to the lexical analyser. Interspersed in the definition of the language may be calls to the error message routines and semantic routines. It is through these calls that the front end is constructed. The parameters at the beginning of the parser, lexical analyser and dictionary ensure that the same token has identical attributes in all the routines. The assembler routines link these parameters. Lexical units are defined and recognised in the lexical analyser and used in the parser. Other terminal symbols, e.g. reserved words, operators and special symbols may be defined in the dictionary and then used in particular places in the parser. All terminal symbols are listed at the beginning of the parser specification and given a unique parse number. This parse number is used by the semantic routines to identify these symbols. All non-terminal symbols used in the parser must define some production of the language and must also be listed at the beginning of the parser specification.

The general syntax analyser package is an LL(K) processor and allows look-ahead to any level in the parser. Look-ahead productions may be defined in the parser but this method was unsuccessful because difficulty was experienced when error processing and recovery was attempted. Look-ahead for terminal

symbols only was used.

There are various options made available by the parser to the compiler writer. One can mark the beginning of a production, inhibit stacking of certain symbols or reserved words, force a particular amount of unstacking say at the end of a production, and inhibit ambiguity checking.

ADA proved to be an ambiguous language and restructuring of productions, look-ahead in productions, semantic analysis and the use of the ambiguity stop had to be used in the specification of ADA for use with GSA 1100.

### 2.1.3. DICTIONARY.

The dictionary is a table containing all key words, symbols and operators. Through the dictionary one can give terminal symbols various attributes that may be used during semantics, e.g. giving order to the boolean values true, false. This routine sorts these terminal symbols into groups of characters for identification. It is important that the parse numbers of these key words is the same in all routines.

### 2.1.4. ERROR MESSAGE TABLE.

GSA 1100 gives the user an efficient way of constructing error messages and recovering to various parts in the syntax. A language and associated construction routines are provided to permit the production of meaningful error messages. As this software is part of the syntax analyser package, and not significant to the research being undertaken, it is not described. Again there is an assembler routine that interfaces with this

routine and the other table generating routines or user routines. It transfers the error token around and recovers to the correct places in the parser. An error recovery table is set up by the parser if there are any calls to the error message routine in the parser specification, and then the relevant error message is built as necessary. The call to the error message routine is done by referring to an error message number, or, in user-defined procedures, referencing the entry point of the error routine with the relevant error message number.

It can be seen that the above routines are very powerful and with their ancillary routines make an efficient tool in the designing and implementation of a compiler.

## 2.2. ANCILLARY ROUTINES.

Several other routines are required to interface with the table generating routines of the general syntax analyser package. These were largely supplied by UNIVAC and are responsible for initialisation of the system, searching and maintaining the stacks and interpreting the tabular output. A brief description of the most relevant one, token, follows.

### 2.2.1. TOKEN.

This routine is the most important of the supporting routines as it directs flow through all the routines. It is a user-written interface between the parser and lexical analyser and generates a token entry (detailed in CHAPTER 5) that is acceptable to the parser and any semantic routines that may reference it. While doing this it -

- 1) recognises key and reserved words by accessing the dictionary (the tables produced by the keyword builder).
- 2) maintains a dictionary of items recognised and so builds a symbol table.
- 3) converts and handles any special class of lexical units that map onto different token units. e.g. Certain words carry semantic information but the parser only needs these words to be recognised as identifiers.

Input to the token routine comes from the GSA 1100 lexical analyser. Output from the token routine consists of a single token entry that is placed in the input buffer of the parser.

The module containing the token routine has been extended to contain all the semantic checking. (See section 2.3 and CHAPTER 5).

The structure of GSA 1100 is modular, thus enabling the user to produce one error message table and one dictionary which may be used with all future releases of a language. Every time the language specification is changed, only the parser and lexical analyser need be updated. Emphasis can therefore be placed on implementing the language and not necessarily on error message production and recovery, which, once developed, remains static.

Figure 2-2 gives an overview of the way which the GSA 1100 table generating routines and interfacing routines work together to produce meaningful output.

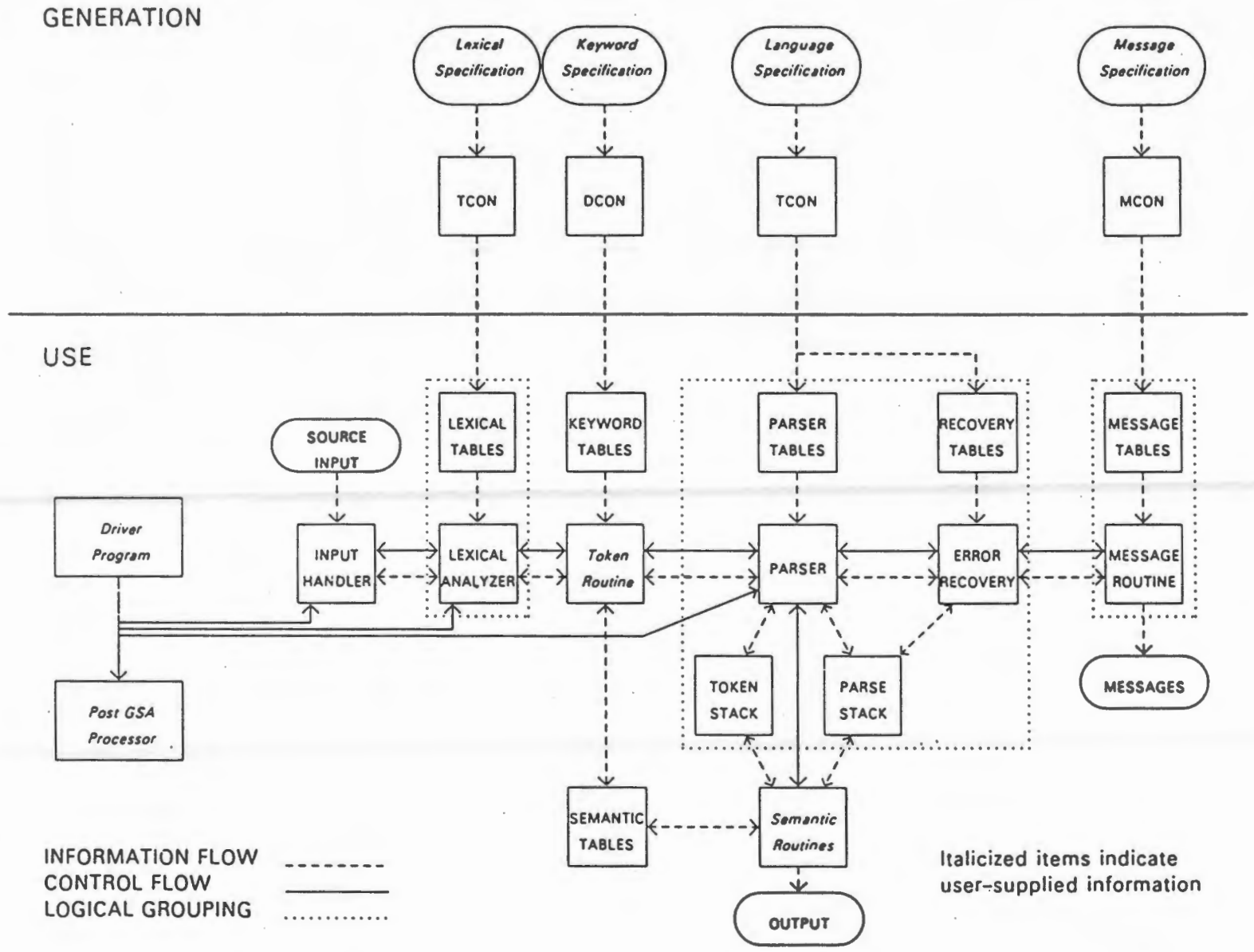


Figure 2-2. GSA 1100 SYSTEM OVERVIEW. [ Sperry Univac 2 ]

### 2.3. STATIC SEMANTIC ANALYSIS.

Static semantic analysis is done through several procedures which reside in the module token. This was done because the token routine that interfaces with the GSA table generating routines produces the symbol table, token stack and various other arrays and stacks that are used in the semantic analysis. All the calls to these semantic routines are either imported from the parser or generated by other semantic checking routines.

Most of the static semantic checking is type checking which is achieved by using the information in the symbol table together with the token stack. Intermediate results are stored on the token stack and where static expressions are evaluated, the results are stored in an array of values and a pointer is set up from the symbol table entry to the expression value.

Overloading within numeric types and the converting of strings of numeric digits to their numeric value is catered for. Where possible, range checking in the form of type compatibility and non-null ranges is done, and each array subscript is tested for compatibility between it and the allowed subscript type for that array.

For a detailed description of these procedures, see CHAPTER 5.

The intermediate code being produced is DIANA, and where possible, the DIANA constructs include semantic information e.g. numeric values, intermediate types etc. Thus calls to the DIANA producing routine are made from the semantic analysis procedures where applicable.

The token module is therefore the pivot of the whole front end processor.

## 2.4. PRODUCTION OF DIANA.

Two of the original design goals of DIANA were:

- 1) " DIANA is based on the Formal Definition of ADA " and
  - 2) " in DIANA, there is a single definition of each ADA entity "
- [Goos 1981].

These two goals conformed to the method of construction of the UCT-front end and resulted in calls to the DIANA producing routine primarily being made from the parser. As each construct is identified in the parser, the relevant DIANA code is produced with any associated semantics from the symbol table.

It is against the batch processing philosophy to produce code for an incorrect program. Thus, once a compilation error is found, no more DIANA is produced. This allows the production of DIANA to assume a syntactically correct ADA program as input.

Another of the design goals was uniformity within the DIANA constructs. Consequently different DIANA structures may have similar configurations which rely on scope and context sensitive information for their construction. Scope problems were overcome by using several stacks keeping track of the position of certain nodes in the DIANA tree, block numbers, beginning of lists, etc. Once a construct is complete, it is output to a file.

As each DIANA node is generated, the DIANA routine attempts to complete it. Every DIANA node has the same basic structure and may be interpreted through a unique identifying number. By using this number and correlating it with the kind of node, one can determine what the subsequent attributes mean.

The production of this intermediate code is performed in the module DIANA which is written in PLUS. For a detailed description see CHAPTER 6.

## 3. CHAPTER 3. BACKGROUND INFORMATION.

## 3.1. ADA.

## 3.1.1. HISTORY.

In the early 1970's the United States Department of Defence decided that it was necessary to do something about the rising software costs; the most important of which was that of the embedded system sector (approximately \$1680 million). To counter this it was decided that a single standard was needed for this sector.

A set of requirements was developed and refined in a series of documents.

DATE	DOCUMENT	PURPOSE
1975	Strawman	Informal comments
1975	Woodenman	Informal comments
Jan 1976	Tinman	Proposed language
Jan 1977	Ironman	Informal comments
April 1977		Call for tender
July 1977	Revised Ironman	
	PHASE 1 GREEN RED BLUE YELLOW	
	Steelman	
April 1978	PHASE 2 GREEN RED	
	Preliminary ADA	
May 1979	PHASE 3 GREEN	
July 1980	ADA-80	
July 1982	Draft ANSI standard	

The call to tender for the language in April 1977 received seventeen replies. In August 1977, when Phase 1 began, four of these were chosen. Each was designated a colour code: CII-HONEYWELL-BULL (green), INTERMETRICS (red), SOFTECH (blue), and SRI INTERNATIONAL (yellow). In May 1979 the green language was accepted. This was designated Preliminary ADA.

After much criticism and experimentation, changes were made to preliminary ADA [ Preliminary 1979 ] resulting in ADA-80 [ Military 1982 ]. ADA-80 overcame many of the ambiguities and semantic problems that existed in the preliminary language.

ADA-80 was subsequently reviewed and modified in the light of implementation difficulties and a draft standard document was released [ Reference 1982 ]. This document became available to the public in October 1982 and appears to have resolved and clarified many of the hazy areas of the previous two releases of the language.

ADA compilers fulfilling the validation requirement of the Department of Defence [ Goodenough 1981 ] are being produced internationally. Of necessity these compilers' abilities include the specified tasking capabilities of the language. All leading manufacturers and several large institutions are currently supporting, sponsoring and implementing ADA. This should result in ADA becoming the real-time standard of the 80's.

An important complementary activity was the development of a parallel series of documents entitled SANDMAN, PEBBLEMAN and STONEMAN as the ADA support environment [ Stenning 1981, Wolfe 1981 ]. This enhances program portability and allows programmers' proficiency in one standard language.

## 3.1.2. THE LANGUAGE.

ADA is the programming language designed according to the STEELMAN requirements of the United States Department of Defence [ Steelman 1978 ]. ADA was named after Ada Lovelace who was a leading computer pioneer of the nineteenth century, a colleague of Charles Babbage, and the daughter of Lord Byron. It was designed at CII-HONEYWELL-BULL by a Paris based team led by Jean Icbiah.

The main basis of ADA is the programming language PASCAL and its later derivatives. A goal in the design of ADA was to retain the Pascal spirit of simplicity, but not necessarily the form of each Pascal feature.

ADA possesses the data abstraction features of modern programming languages, and as such, has the ability to define types and subprograms with the conventional control structures serving the need for modularity. An integral part of the language is real-time programming with the ability to model parallel processing and cater for exception handling. Included are systems program applications requiring access to system dependent parameters and precise control over the representation of data.

The three design criteria for ADA were

- 1) the importance of program reliability and maintenance.
- 2) the fact that programming is a human activity.
- 3) programs are required to be efficient.

The designers endeavoured to make it easy for programmers to write good, structured programs using ADA.

An attempt was made to keep the language as contained as possible. This gave rise to a language with a minimum number of underlying concepts put together in a consistent and ordered manner.

A program in ADA is a sequence of higher level program units.

Program units may be subprograms, package modules or task modules. Separate compilation of program units allows a program to be designed, written and tested in independent parts.

A subprogram may be either a procedure or function subprogram, and is the basic unit for expressing an algorithm. A procedure subprogram performs a sequence of desired actions, whereas a function subprogram computes a value. A package module defines a collection of logically related units. Portions of packages can be hidden from the user, allowing access only to the logical properties of the module. A package module may be a collection of related subprograms.

A task module also defines collections of logically related units but with the additional capabilities of parallel processing. Tasks may be implemented on multiple processors, or interleaved on a single processor.

### 3.1.3. VERSIONS OF ADA.

Preliminary ADA was released in 1979 [ Preliminary 1979 ] and had some very interesting features, e.g. dynamic arrays, initial default values for parameters, the idea of separate compilation units, etc. It left many areas unexplained and ambiguous. This version was studied and discussed by people throughout the world and when ADA-80 was released there were significant changes. For a detailed list see [ Winkler 1981 ].

One noticeable difference between the two documents was that the ADA-80 specification was more detailed and more precise than Preliminary ADA. In several instances restrictions and details were specified where previously no elaboration had existed.

The UCT-front end is based on the ADA-80 definition of the

language [ Military 1980 ] where there are still several flaws with which implementors have to contend. The draft standard document for ADA [ Reference 1982 ] overcomes several of the problems present in ADA-80. Below are a few of the syntactical changes between ADA-80 and "final" ADA. Only changes relevant to the UCT-front end are mentioned.

The new manual explicitly details the semantics and rules of the language including notes to the implementor about the interpretation of the explanations.

The syntactical differences included in this implementation are:

- 1) NAME has been restructured. The UCT-front end has taken the final ADA definition of SELECTOR.
- 2) The relation part of an EXPRESSION has been simplified and this simplification has been implemented.
- 3) Wherever an EXCEPTION\_HANDLER appeared in ADA-80, it was optional, even if the reserved word EXCEPTION was present. This has been changed. If the word EXCEPTION is used, it must be followed by at least one EXCEPTION\_HANDLER. This has also been implemented in the UCT-front end.

Other changes are:

- 1) DECLARATION has been changed to overcome an ambiguity that existed in the syntax for DECLARATIVE\_PART. Several productions that may have appeared further down the parse tree have been moved so that the correct path can be identified as early as possible. Consequently DECLARATIVE\_PART and all the productions it calls have been redefined.
- 2) An OBJECT\_DECLARATION can no longer contain an unconstrained ARRAY\_TYPE\_DEFINITION. This infers that all unconstrained array types must be explicitly declared as types.
- 3) The syntax for TYPE\_DECLARATION has been altered to include a

PRIVATE\_TYPE\_DECLARATION. This has no actual significance in the structure of a TYPE\_DECLARATION as the PRIVATE\_TYPE\_DEFINITION has been taken out of TYPE\_DEFINITION causing no difference to the resulting syntax.

- 4) RANGE has been changed to include the range attribute.
- 5) In several places where previously a SUBTYPE\_INDICATION was specified, ADA now only requires a TYPE\_MARK. Also where NAME was specified, ADA now requires only a SIMPLE\_NAME which is an IDENTIFIER. These two simplifications make the whole syntax less ambiguous and easier to implement.
- 6) The components of a record may no longer be declarations of arrays but just variables of an existing user-defined or standard type. Also, a COMPONENT\_LIST must now have at least one component, even if it is the NULL component.
- 7) ALLOCATOR has been simplified.
- 8) One of the major differences is that of a parameterless function call. In ADA-80, an empty pair of parentheses had to follow the call. This is no longer required.

The remaining syntactic changes refer to packages, tasks, generics and library units, all of which have not been implemented, so are not detailed here. They may be found in [ Harrison 1982 ].

The changes made to "final" ADA are mainly for the better as they bring ADA in line with other current languages. Although programmers are accustomed to the change in FUNCTION\_CALL, this will continue to cause further ambiguities for the implementor.

## 3.1.4. RESTRICTIONS.

As the implementation of a compiler for the complete ADA specification would involve tens of man years work, only a subset could be implemented. The syntax checker has been constructed for the complete ADA definition ( except the RENAMES clause ), but only a subset has been extended further.

The following rationale was used to extract the subset. A program in ADA is a sequence of higher level program units. Program units may be subprograms, package modules or task modules.

The basic unit for expressing an algorithm is a subprogram which may be split into procedural subprograms and function subprograms. A package module defines a collection of logically related entities. A task module is basically a package module with additional capabilities for parallel processing. Since modules require the basic constructs of a subprogram, it was decided to implement only subprograms. This subset is equivalent to PASCAL. Just as PASCAL can easily be extended to concurrent PASCAL, the selected subset can be extended to encompass the whole of ADA using GSA 1100.

In particular the following were not implemented: pragmas, modules, visibility rules, tasks, program structures and compilation issues, generic program units, representation specifications, implementation dependent features and the RENAMES clause.

Pragmas are messages to the compiler. They were excluded because insufficient information was known about GSA 1100 to be able to implement pragmas successfully.

Task modules and package modules were excluded since they are not really applicable in the present environment. Generally, students

do not work in groups on a single package, and therefore the need for packages is small. On a single processor system tasks are usually only useful for simulation, and therefore they were excluded.

Since the visibility rules are intended for modules, they were removed from the workable subset. Compilation issues are used mainly for packages in program libraries, and are not relevant to a subset excluding separate compilation and modules.

Other parts were originally excluded from the subset, but were later inserted as their meaning became clearer and they were seen to be part of the basic subset. This subset has been implemented as far as the production of DIANA, that is syntax checking, static semantic checking and the production of DIANA. As the remainder of ADA has been implemented in the syntax checker, this implementation can be extended through the semantic phase to the production of DIANA to give a complete front end for ADA.

## 3.2. DIANA.

### 3.2.1. INTRODUCTION.

DIANA, a Descriptive Intermediate Attribute Notation for ADA is based on two other intermediate codes - TCOL and AIDA. DIANA is expected to become the standard intermediate language for ADA implementations, and so we decided to use it as the end point for the UCT-front end.

The design of the language, as stated in the DIANA manual, [ Goos 1981 ] is based on the following principles:

- 1) It is representation independent.

- 2) It is based on the formal definition of ADA.
- 3) Regularity is a principal characteristic of DIANA.
- 4) It must be efficiently implementable.
- 5) Consideration of the kinds of processing to be done is paramount.
- 6) In DIANA, there is a single definition of each ADA entity.
- 7) DIANA must respond to the issues posed by ADA's separate compilation facility.
- 8) There must be at least one form of DIANA representation that can be communicated between computing systems.

These design goals give rise to an intermediate language that has a syntax similar to ADA produces uniform structures for the different productions and is easy to implement as the user sees fit. The resulting intermediate form of the program resembles the original program in such a way that there should be no problem of regenerating the source, or of designing tools which use the "DIANA" program.

Since the original release of DIANA (March 1981) several problems have been encountered. These have been studied and a new draft release of DIANA (October 1982) [ Evans 1982 ] attempts to overcome these. DIANA-82 is compatible with the draft proposed ANSI standard for ADA.

As the UCT-front end is based on ADA-80, it used DIANA-81.

### 3.2.2. THE LANGUAGE.

An instance of the DIANA form of an ADA program is an attributed tree. The structure of the tree is that of the abstract syntax tree defined in the ADA formal definition. Attributes of the nodes of this tree encode the results of semantic analysis.

The DIANA tree is structured according to the syntax of ADA. Each ADA production has a corresponding DIANA production giving rise to a set of nodes for that production. Any static semantic analysis is held in these nodes through the semantic attributes of the node. If that node relies on system supplied information, there will be attributes for this information specified in this node. In this manner, the DIANA tree contains the results of the syntax analysis, static semantic analysis and system supplied information.

Every node has the attribute LX-COMMENTS and most have the attribute LX-SRCPOS. The latter is the source position of the symbol and these two attributes are for use in regenerating the source code.

ADA SYNTAX : ASSIGNMENT\_STM ::= VAR\_NAME := EXPRESSION

DIANA SYNTAX : ASSIGN => as-name : NAME  
as-exp : EXP

ASSIGN => lx-srcpos : SOURCE POSITION  
lx-comments : COMMENTS .

EXAMPLE A := 3

The DIANA tree to be produced would be

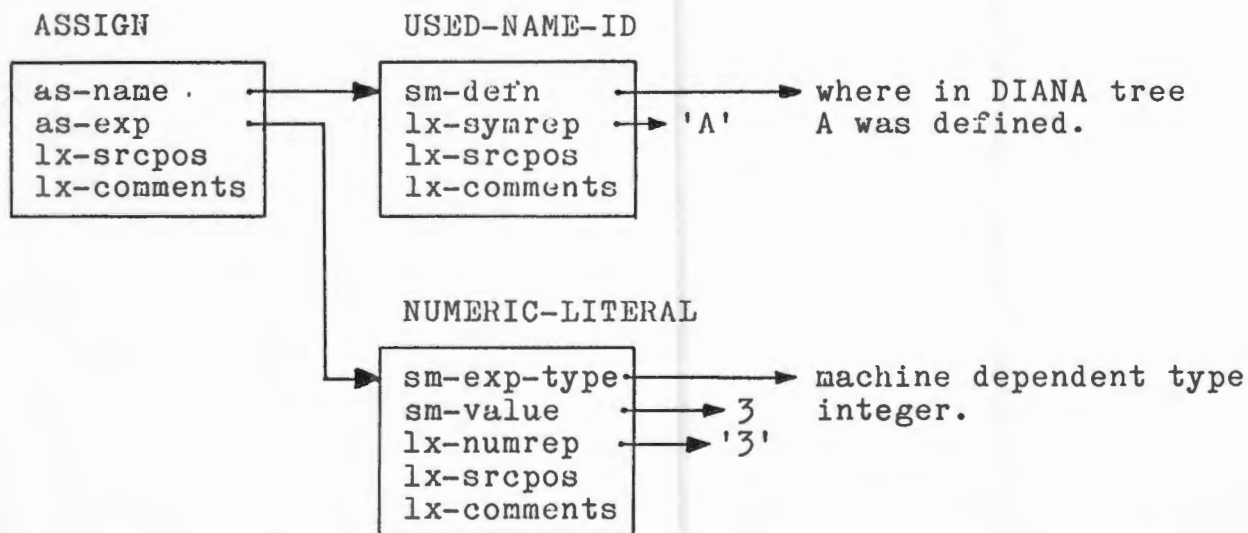


Figure 3-1. EXAMPLE OF AN ADA STATEMENT AND ITS DIANA TREE.

Figure 3-1 is an example of an ADA statement and its equivalent DIANA tree. The section of tree produced would be linked to the other statements in the block of statements, or to the block header.

DIANA has proved easy to implement, especially due to the almost one to one mapping between each ADA production and each DIANA tree.

## 4. CHAPTER 4. THE SYNTAX CHECKER.

Four modules constitute the Syntax Checker. They are the lexical analyser, parser, dictionary and error message table. The pivot of these is the parser. It uses information from the other three modules to do the syntax checking specified.

CHAPTER 2 briefly describes the function of these modules and this chapter will only discuss how ADA has been implemented using the facilities of the GSA 1100 parser.

The parser is a top-down, look-ahead parser that uses the tables built from the language specifications to control parsing. It can be envisaged as an interpreter that has a repertoire consisting of recognition, control and action instructions.

The token routine passes input recognised in the lexical analyser to the parser through a "window". Depending on the input and the present state of the parser, the appropriate action is taken. As input is recognised it is placed on the token stack available for semantic routines.

Look-ahead is used to resolve ambiguities by specifying look-ahead at terminal symbols in the input stream. It is specified and interpreted as follows:

```
<STATEMENT> ::= <NAME> ( '(' <PARAMETER_LIST> ')' )
                ?[ ':=' : ':=' <EXPRESSION> ]? ';' .
```

This is read as :

- 1) Accept a <NAME> .
- 2) If the next token is a '(' accept '(' <PARAMETER\_LIST> ')'
- 3) Look ahead at the next symbol. If it is a ':=' then accept ':=' <EXPRESSION> and proceed to 4). If it is not ':=' then proceed to 4.
- 4) Accept a ';' .

Sets are used to specify the recognition of any set of terminal symbols e.g. all the names that could be parsed as identifiers.

Interspersed between the productions making up the ADA language are calls to the error message routines indicating which error message to print and what recovery is needed, calls to the token module routines controlling semantic checking, and calls to the DIANA module directing the production of the DIANA code.

In implementing a syntax checker for ADA, the semantics of the language were ignored and only the syntax considered. This was found to be ambiguous and overspecified. Due to the structure of the GSA 1100 parser, these problems became the most difficult to solve of those encountered.

Ambiguity means that at a certain point in the syntax, it is not clear from the input item which path in the syntax is being followed. Overspecification is when alternate paths in the syntax cannot be distinguished from one another. Overspecification is the worst case of ambiguity.

#### 4.1. AMBIGUITY RESOLUTION.

An example of the kind of problem encountered can be seen in the condensed version of the following two productions (ADA manual sections 3.7.3, 4.3 [ Military 1980 ] )

```
<CHOICE> ::= <SIMPLE_EXPRESSION> \ <DISCRETE_RANGE> \ OTHERS .
<COMPONENT_ASSOCIATION> ::= ( <CHOICE> ( '|' <CHOICE> ) ... => )
                                <EXPRESSION> .
```

Suppose a token is recognised as:

- 1) IDENTIFIER. This may be parsed as a <DISCRETE\_RANGE>, <SIMPLE\_EXPRESSION> or <EXPRESSION> .
- 2) NAME. The same three options exist.

3) SIMPLE\_EXPRESSION. The same three options exist.

It can be seen that even after recognising a complex structure such as <SIMPLE\_EXPRESSION>, there is no way of knowing which path to follow.

This ambiguity does not exist only if an identifier is found, but also if a number, CHARACTER\_STRING, ALLOCATOR or left parenthesis is found. For example, say a '(' is found. The next token may be an expression of any complexity followed by the closing parenthesis ')'. The sequence "( EXPRESSION )" can be a parenthesised EXPRESSION or an AGGREGATE. Both of these occur as a PRIMARY. This may be recognised as <SIMPLE\_EXPRESSION>, <DISCRETE\_RANGE> or an <EXPRESSION>.

Ambiguities occur frequently, and where one occurs, it may involve many productions due to the design goal of a minimum number of underlying concepts. They are resolved in one of three ways, restructuring of the productions, through the use of look-ahead and with the aid of semantic information.

Restructuring involves combining the ambiguous parts of each production to create a new production used by the previous two ambiguous productions.

This technique may be used in conjunction with the look-ahead technique to resolve the ambiguity completely. (see example, figure 4-1). Look-ahead is done only for terminal symbols. This may require specifying several terminal symbols before reaching one that will identify which path to follow.

In the following example, restructuring of the syntax and look-ahead have been used to resolve this simple ambiguity.

Figure 4-1. EXAMPLE SHOWING AMBIGUITY RESOLUTION BY RESTRUCTURING

```

<DECLARATION> ::= [ <TYPE_DECLARATION> \ ....
                  \ <OBJECT_DECLARATION>
                  \ <EXCEPTION_DECLARATION>
                  \ <NUMBER_DECLARATION> ] .

```

```

<OBJECT_DECLARATION> ::= <IDENTIFIER_LIST> ':' ( CONSTANT )
                        <SUBTYPE_INDICATION>
                        ( ':' <EXPRESSION> ) ';' .

```

```

<NUMBER_DECLARATION> ::= <IDENTIFIER_LIST> ':' CONSTANT
                        ':' <EXPRESSION> ';' .

```

```

<EXCEPTION_DECLARATION> ::= <IDENTIFIER_LIST> ':' EXCEPTION ';' .

```

The ambiguity in the last three productions is clear, and can be resolved only after the colon. An ambiguity still exists between <OBJECT\_DECLARATION> and <NUMBER\_DECLARATION>, so look-ahead was used to look for the assignment sign which identifies the number declaration. The resulting declarations are:

```

<DECLARATION> ::= [ <TYPE_DECLARATION> \ ....
                  \ <IDENTIFIER_LIST> ':'
                  [ <EXCEPTION_DECLARATION> \
                    ?[ CONSTANT ':' : <NUMBER_DECLARATION> -> Q ]?
                  <OBJECT_DECLARATION> ] ] ^ Q .

```

```

<OBJECT_DECLARATION> ::= ( CONSTANT )
                        <SUBTYPE_INDICATION>
                        ( ':' <EXPRESSION> ) ';' .

```

```

<NUMBER_DECLARATION> ::= CONSTANT
                        ':' <EXPRESSION> ';' .

```

`<EXCEPTION_DECLARATION> ::= EXCEPTION ';' .`

A common cause of ambiguities was the use of NAME. NAME starts with an IDENTIFIER, where an IDENTIFIER is any sequence of letters, digits and underscores. Often the type of IDENTIFIER used was context dependent and this was used in solving the ambiguities by semantics.

Every IDENTIFIER used in a DECLARATION was flagged as having a particular characteristic, e.g. variable identifiers, type identifiers, constant identifiers etc. Where it could be determined that a particular context permitted only a particular type of identifier, this is all that was allowed. This idea worked very well but became too restrictive and began trapping semantic errors before the syntax had been checked. As a result, semantic checking of this sort is used with care.

Another kind of semantic checking used is where the path to take is determined by characteristics of the token just parsed. This method solved the terrible ambiguity of NAME. An array reference has a subscript list which is all part of NAME and a PROCEDURE\_CALL may have a PARAMETER\_LIST. With a PROCEDURE\_CALL only the procedure NAME is parsed as part of NAME.

Oversimplified, both start with a NAME optionally followed by '(' IDENTIFIER. The ambiguity is when the open parenthesis is recognised. Is it part of NAME or part of a PARAMETER\_LIST?. This was resolved by checking to see if the NAME before it is an array NAME. If so, the whole construct to the next close parenthesis is parsed as NAME otherwise the NAME ends before the open parenthesis.

In the article "Ada Syntax Diagrams for Top-Down Analysis" [Bonet], R. BONET et AL discuss how they reduced the ADA syntax to be LL(1). The methods used are very similar to those used by the

UCT implementation. Their restructuring appears to be more ordered and they make greater use of semantic information to achieve their goal.

The UCT syntax checker implements the whole of ADA (excluding the RENAMES clause and separate compilation) and does full error checking with recovery where possible. There are areas which could be improved and the ambiguities resolved in a more efficient manner. This was not done as the next phase in the development of the front end appeared to be more challenging.

## 5. CHAPTER 5. THE TOKEN MODULE.

Static semantic checking constitutes all checking that can be done at compile time. The procedures that do this checking are all found in MODULE TOKEN and use the tokens on the token stack together with their symbol table information. Results of this checking may be temporarily stored on the token stack for immediate re-use or stored in various other stacks and arrays linked to the symbol table.

## 5.1. STRUCTURE OF THE SYMBOL TABLE.

The symbol table is a structure consisting of linked lists of nodes, where each list header is accessible through a hashing function. Each node is a symbol table entry and may be referenced through a pointer to the symbol table. Symbol table entries take the form of a PLUS (Programming Language for Univac Systems) record, each field holding different semantic information or references to other structures holding information.

A symbol table entry occupies seven words of storage. These are divided into the fields illustrated in figure 5-1 which hold the following information:

## 5.1.1. STRUCTURE.

The structure of a token is the number of characters and words making up the token, its lexical level, block number and lexical type. Once a token has been identified by the lexical analyser, these values are known or calculated and stored for further reference.

BITS WORD	0	3	6	9	12	15	18	21	24	27	30	33	36
1	PARSE TYPE			NUM. CHARS			LEX TYPE	CLASS VALUE			NUM. WORDS		
2	CLASS			LEX LEVEL			NEW ENTRY			BLOCK NUM.			
3	POINTER: NEXT-DICT						POINTER: ITEM-LINK						
4	BASE TYPE		COMP TYPE		PARSE DEF			IND EX	I N C O M P L	D I M E N S I O N S	D I M L I N K	D E R I V E D	D Y N A M I C
5	POINTER: TYPE-PTR						POINTER: NEXT-VAR						
6	SM-DEFN						BSETPE						
7	TPE-STRUCT						UNUSED						

WORDS 1-5 CONTAIN INFORMATION FOR STATIC SEMANTIC CHECKING.

WORDS 6-7 CONTAIN SEMANTIC INFORMATION USED BY THE DIANA MODULE.

Figure 5-1. STRUCTURE OF A SYMBOL TABLE ENTRY.

#### 5.1.2. SEMANTIC INFORMATION.

Semantic information is divided into two categories, the information used in the static semantic checking and the information necessary to complete the semantic pointers of DIANA nodes.

Information used for static semantics includes:

1) PARSE-TYPE. A number used to identify a particular group of terminal symbols used in the parser. e.g. If an identifier is

defining a new type, the identifier will have a parse-type indicating a type-name.

2) PARSE-DEF. A number used to identify the type definition being defined. e.g. Record, array.

3) CLASS-VALUE. Each predefined type has a unique type number for identification. Each user-defined type must also be given a unique type number. These are stored in the CLASS-VALUE field. This field is also used to indicate whether a variable is defined to be a constant.

4) CLASS. Within each CLASS-VALUE there may be several unique groupings to be identified. e.g. The type BOOLEAN has two values, TRUE and FALSE. They are defined to have a particular order and this is held in this field.

The values for the PARSE-TYPE, CLASS-VALUE, and CLASS may be pre-set in the dictionary.

5) BASE-TYPE. The type of the first declaration in a string of declarations constitutes the base-type of a variable. Overloading may require that the base-type of a definition be used. This could be found by backtracking through the connecting pointers. The structure of the first declaration is useful in determining properties of a variable and this could also be obtained through backtracking. If only backtracking were used the presence or absence of the certain properties would be recalculated each time a particular variable was used. Therefore two extra fields holding the BASE-TYPE and component type (COMP-TYPE) of declarations are included in each entry. The component type is only meaningful when a complex structure e.g. array or record is being defined.

6) Various other fields indicate whether or not a type declaration is complete or derived and if the entry in the symbol

table is a new entry. If it is, certain semantic information must be supplied and the new entry information is used by a routine called from the parser.

7) In each entry for an array variable is a field holding the number of dimensions of the array.

### 5.1.3. POINTERS.

Several tokens may constitute a complex structure, e.g. record. The elements of this structure are linked together by Pointers. Pointers are also used to link variables to their type definitions.

Numeric values and the results of expression evaluation are stored in arrays. The value of the indices to these various arrays are stored in particular fields in each entry. Arrays may have indices of different types. These types are stored in a type array and the position in this array is specific to a particular definition, and is stored in the symbol table entry.

### 5.1.4. REPRESENTATION.

At the end of every entry in the symbol table, the ASCII representation of the token is stored. The number of words needed to do this is calculated by the routine token and is stored in the symbol table for future use.

DIANA symbol table information is completed and referenced in the DIANA module.

Figure 5-2 shows the basic symbol table entry showing linkage for the following extract of ADA.

```

TYPE ONE IS RECORD
  D,E : INTEGER ;
  F   : FLOAT   ;
END RECORD ;
TWO : ONE ;
    
```

Figure 5-2. SYMBOL TABLE LINKAGE.

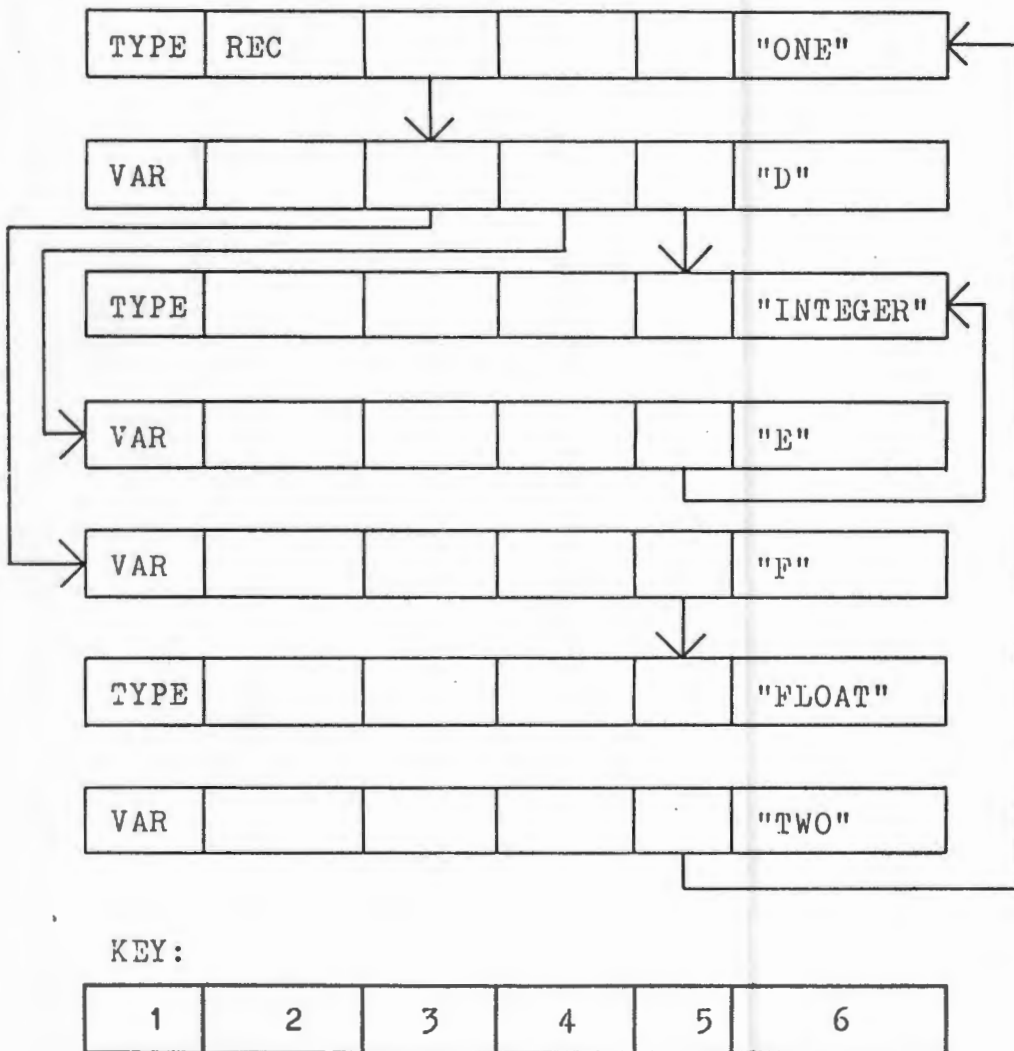


figure 5-2. CONTINUED.

- 1: PARSE-TYPE of entry.
- 2: PARSE-DEF of entry.
- 3: Pointer linking the record definition to the first element of every list in the component list.
- 4: Pointer linking elements making up each list in the component list.
- 5: Type pointer linking the variable to its type definition.
- 6: Representation.

## 5.2. FUNCTION OF THE TOKEN MODULE.

### 5.2.1. PRE-PARSER PROCESSING.

Lexical units are constructed by the lexical analyser and passed to the token procedure of the token module with the number of characters constituting the unit. The token procedure uses the number of characters to calculate the number of words the symbol will occupy, and the hashing function value needed to access the symbol table. It then searches through the symbol table, using the hashing information to establish if the symbol has been entered before. If not, a new entry is created. This new entry contains the symbol representation, parse and lexical types (which remain the same until the parse-type is modified during post-parser processing), number of characters and number of words making up the symbol, and a flag which is set to indicate that a new entry has been created. Once the entry has been created a check is made to see if it is an identifier or a special symbol. If it is an identifier, it could be a reserved word which has

special properties defined in the dictionary.

The results of this check, direct the token procedure to return either the token to be passed to the parser, or to search the dictionary for additional information before returning the token. The dictionary supplies the token entry of specified tokens with the parse-type, class and class-value. If the tokens are not specified in the dictionary, these fields will be completed during post-parser processing.

The search of the symbol table will often result in a match. This may or may not be a legal occurrence of the token so the token, with its semantic information is returned to the parser for post-parser processing.

#### 5.2.2. POST-PARSER PROCESSING.

##### DECLARATIONS.

CHAPTER 4 refers to using semantic information to solve ambiguities in the ADA language. The semantic information used is that obtained from the declarations of the variables and their types.

ADA defines several forms of named entities, the most important being TYPE, NUMBER and OBJECT declarations. Each TYPE\_DECLARATION defines a particular structure e.g. record, array, access type, etc. Each DECLARATION creates an IDENTIFIER with a special function and properties, and perhaps structure. This information must be encoded into the symbol table, and this is done through calls from the parser.

Every declarative production initiates a call to a "declaration" procedure in the token module. This procedure accepts the token created or found during pre-parser processing. If the new entry

flag is true, it is an IDENTIFIER that has not been declared before, and the PARSE-TYPE field can be completed without any further processing. This routine may encounter identifiers that have been declared before, either at an earlier lexical level or in error. If the new entry flag is false, the lexical level of the symbol table entry and the current lexical level are compared. An error has occurred if the levels are the same, but if they are different a new entry is created for the IDENTIFIER.

The entry for the IDENTIFIER is marked with a pointer to enable it to be completed as more semantic information pertaining to its declaration becomes known. A call from the parser would indicate the type of structure being defined and this is entered in the PARSE-def field of the entry.

The pointers are all created from one entry to another as a result of calls from the parser.

Every time the symbol is referenced in a production that is not a declaration, the above information is available for use in semantic checking. Figure 5-2 indicates that this semantic information constitutes the most important fields in the symbol table. The remaining fields are completed during the parsing of the body of the program.

#### SCOPE.

Scope is an important factor in any computer language. Where the same variable is being defined twice, the symbol table relies on lexical levels to determine whether the definition is a legal redefinition of an inner level, or an error.

The token module houses routines that increment and decrement block numbers and lexical levels at the appropriate times.

## INITIALISATION.

The token routine performs certain initialising routines.

It assigns initial values and properties to the predefined types e.g. INTEGER, FLOAT, BOOLEAN etc. They have been defined to have a lexical level of zero and may be redefined by the user at a higher lexical level. Thus each of these types is an entry in the symbol table and may be referred to just as any other entry. Any predefined entry that may be redefined must be initialised in token and not in the dictionary because the dictionary only allows one definition of each key word, whereas the token routine is structured to allow redefinition of identifiers at different lexical levels.

Flags, arrays and counters are explicitly initialised when they are declared, or in an initialising procedure.

## STRUCTURES AND STACKS USED BY THE TOKEN MODULE.

1) The token stack is a stack created by the parser. Tokens of a production are loaded on to this stack and are accessed by the semantic routines. Each token on the stack may be linked to its symbol table entry to enable the semantic routines to access the information of the token. Using this stack, each token entry can be traversed and evaluated or inspected. A pointer indicates the current position in the stack and may be moved along the stack for whatever processing is required. The implementor may add entries to this stack for semantic use and it is through this stack and its pointers that most of the semantic checking is performed and transferred to the symbol table.

2) Each declaration of an array may have several indices, each of which may have different types. These types are stored in an array. The number of dimensions of the array declaration is

stored in the symbol table entry for that array and its number of indices is therefore known. The index to the type-array which refers to the type of the first index is stored in the symbol table and by using the number of dimensions field in the symbol table the others may be located for subscript checking. Used in conjunction with this type-array is a bound-array. It is used not only for array variables but for any variables with an associated range. If the range is static, the bounds of the range are stored in this bound-array where they are accessed for range checking. A similar mechanism to that of the type-array index is used for the bound-array.

3) Numeric strings are converted to their equivalent numeric value and stored in one of two arrays. There is an array for storing reals and one for integers. Once a string has been converted, its value is linked to the symbol table entry. The next identical string recognised need not be converted to its numeric equivalent as this has been stored.

4) Type compatibility is a major part of the static semantic checking. A complex structure may not necessarily be used in its most complex form, thus type checking does not involve using the base type of a structure. Often the type of an intermediate structure is temporarily required and after some processing, the type of the next stage of the structure is required. An array of pointers is used to keep particular positions in the structure for further reference. This array works as a stack, increasing and decreasing with complexity of the structure.

## EXAMPLE.

```

TYPE R IS ARRAY(1 .. 10) OF INTEGER ;
TYPE S IS ARRAY(10.. 20) OF INTEGER ;
TYPE S1 IS ARRAY(10.. 20) OF FLOAT  ;
TYPE U IS ARRAY(1 .. 10) OF S1  ;
I : INTEGER ;
B : S ;
A : R ;
D : U ;
.
.
.
D(B(I))(A(I)) := 3.4

```

Sequence of events.

- 1) Find the type of D( ).
- 2) Check subscript type B(I). This entails leaving the current chain of backtracking to ascertain the type of B( ) and then checking the index I. Once B(I) has been verified against the subscript type of U, control must return to the chain of backtracking which reached D( ) to continue with D( ) ( ). This is done by storing the symbol table entry pointed to by the backtracking pointer in the array of pointers. This method continues until the type of the structure has been reached.

Using the same declarations, and introducing the declaration

```

E : S1 ;
D(3) := E ; is legal.

```

This is an instance when the sequence of backtracking events does not need to go all the way to the base type.

5) Static semantics involve evaluating static expressions. The operands of an expression are not evaluated in a sequential manner. Intermediate results, and types needed for re-use are stored using two stacks, one for the values and one for the types. The stack with the intermediate types also indicates whether the result of an evaluation is a constant. This is used in overloading resolution of arithmetic expressions. Several other stacks are used in particular procedures to perform very specific tasks.

#### OVERLOADING.

The UCT-front end does not yet cater for overloading of the form:

```
TYPE COLOUR IS (RED, GREEN, YELLOW, BLUE, ORANGE) ;
```

```
TYPE RAINBOW IS (RED, ORANGE, GREEN) ;
```

```
VARCOLOUR : COLOUR ;
```

```
VARRAINBOW : RAINBOW ;
```

```
.
```

```
.
```

```
.
```

```
VARCOLOUR := RED ;
```

The last statement may cause an error as the type checking routine may identify RED as being of type RAINBOW instead of COLOUR. This is because the token procedure which sets up the symbol table allows only one occurrence of an identifier at any particular lexical level. In the above example there are two occurrences of RED, ORANGE, and GREEN and an error would have been produced when an attempt was made to parse the type declarations.

This characteristic gave rise to several problems. The first is that the ADA language specification insists that overloading be allowed and the UCT-front end does not allow it. The second is

that records are an integral part of the language. Not to be able to use the same name for fields in two different records seemed a limitation to the user. The later problem was solved in the following manner.

Consider the declarations

```
TYPE A1 IS RECORD          TYPE A2 IS RECORD
  B : INTEGER ;           B : FLOAT ;
END RECORD ;              END RECORD ;
C1 : A1 ;                 C2 : A2 ;
```

Both fields B have different properties and therefore cannot use the same symbol table entry. The symbol table constructing routine allows identifiers with the same name at different lexical levels.

The field B is never referenced without some reference to the variable of type A1 or A2. The components of each record are linked to the record definition (see section 5.1 for structure). Therefore, for the duration of the type declaration, the lexical level is increased to store the components of the record. Once the record has been stored, the lexical level is decreased and all links from the hash table to the record components are removed. This results in the record components only being accessible through their type declarations. Once accessed, the component has all the properties of a normal variable.

The problem then became how to access the component again for use and type checking.

The only two ADA notations for using the components of a record are :

- 1) the point notation e.g. C1.B and
- 2) through the use of an aggregate e.g. C1(B=>2).

In both of these, the C1 is present and can be used to access B by tracing the pointers of the structure.

This is achieved by setting a flag when a '.' or '(' is encountered. Before processing the next token a check is performed to see if the token preceding the '.' or '(' is a record, or is derived from a record. If it is, then the pointers are used to identify the next token and not the hash table method. This solution works well and allows different records to have fields with the same names and differing characteristics. In addition it allows the use of aggregates to set initial default values.

The other form of overloading that was overcome by this implementation was that of overloading within an arithmetic expression. This may occur across types where one operand is a numeric constant and the other is of a type derived from the same type as the numeric constant. This was achieved simply by evaluating both the type and base type of each operand. If one of the operands is a constant, then the base types are used for type compatibility, the result is a "constant", and has the type and base type of the original constant in the operation. This mechanism also caters for arithmetic where operands of different specified types are allowed and the result is also specified e.g. fixed \* float = fixed.

#### EVALUATION OF NUMERIC LITERALS.

Evaluation of numeric literals is not done according to the ADA specification. The UCT-front end evaluates two types of numbers, integer and real. The integer numbers are referred to as ADA UNIVERSAL INTEGER and the real numbers as ADA UNIVERSAL FIXED. All numeric literals are constants. If evaluation of an

expression indicates that a particular constant should have had type float and not fixed, this is rectified.

Every entry in the symbol table is made up of the information relating to the entry followed by the ASCII representation of the token. The tokens are stored four characters to a word and from the token entry one knows how many words make up the token. Using this information each numeric string is converted into its equivalent decimal value. The method used is to decimalise each character by taking the numeric value of the character and multiplying this by the base of the number, creating a cumulative decimal equivalent of the input. e.g. The hexadecimal number 12 would be evaluated as  $1*16+2 = \text{decimal } 18$  and hexadecimal 122 would give  $((1*16+2)*16)+2$ .

A check is made to ensure that no digit greater than the base of the number is used in its specification.

#### TYPE CHECKING.

Type checking covers many areas. There is arithmetic type checking, parameter type checking, subscript type checking, aggregate type checking, range checking, checking the type of initialising values and checking the type mark of a range constraint with the type of the range. This implementation does all of the above except the type checking for aggregates.

In the majority of instances the two operands must have the same type, but for derived types and to resolve overloading, they must have the same base type. The major difficulty in implementing this occurs when the operands are array accesses. The array subscript type has to be matched with the array definition and the array component type with its value. The component type is only determined when the operator separating the operands is reached,

which may be after several other operators have been reached in the subscript.

Evaluating the type of an operand is a recursive procedure, and is repeated to find the type of each array subscript. Several levels of subscripting may exist and these levels are demarcated using the array of pointers.

All instances of type checking use the same procedure to find the type of the operand. The base type of the operand is also determined for use with overloading, derived and access types. Each subprogram declaration is linked in the same manner as a record declaration. Parameter checking therefore involves following the pointers of the structure to find the formal parameters, and matching these with the actual parameters when positional notation is used. Named notation is very similar except the structure may need to be traversed more than once if the named parameters are not in the same order as the formal parameters. The type compatibility between the parameter name and its value is checked, and this indirectly checks that the name is a valid formal parameter of the called procedure. The same method will be used for aggregates.

The ADA language specification explicitly states the typing rules for the arithmetic operators. These were implemented through a case statement, one case for each operator. This catered for all expressions where the type of the operands was universal integer or real, or both operands were of the identical type but not for expressions where the operand types were derived from the universal types. The base type of each operand was introduced to cater for instances of "legal" mixed mode arithmetic.

```
TYPE A IS 1 .. 10;
B : CONSTANT := 3;
C : INTEGER ;
D : A ;
BEGIN
    C := 5 ; --LEGAL
    D := 3 ; --LEGAL
    D := B ; --LEGAL
    D := C ; --ILLEGAL
END ;
```

Figure 5-3. EXAMPLE SHOWING ASSIGNMENT TYPE COMPATIBILITY.

Multiplication and division operators proved interesting. Operations using these operators are allowed between operands of the same base type. Mixed mode arithmetic with these is only allowed between integer and fixed point operands. If the operation between these two types is multiplication, the result is of type fixed. For division, only the left operand type may be fixed and the result is fixed. If both operands are fixed, the result is that of UNIVERSAL FIXED.

These rules have been implemented by flagging operands of type universal as constants, and using this information when storing the eventual type of the result.

Assignment is a special operation included in the case statement for operators. The same rules are applied when checking the type validity of variable initialisation. Two routines are required for these two occurrences of an assignment statement because the routine that determines the type of an operand uses the token stack, and relies on different kinds of tokens to mark the beginning of an operand. These markers are particular to the context.

A + B	OPERANDS ARE A	B
A(1)(B(1)).C-3	OPERANDS ARE A(1)(B(1)).C	3
A : B := NEW C	OPERANDS ARE B	C

Figure 5-4. EXAMPLE OF POSSIBLE OPERANDS.

Range checking requires some type checking and bound checking if the bounds are static expressions. The type checking is of two kinds:

- 1) The type mark of the range constraint must be compatible with the expressions denoting the range.
- 2) The expressions denoting the range must be type compatible.

The type checking routines use the same routine to determine the type of the operands. If none of the type checking fails, range checking is performed and the relevant error messages returned. An important area in this type checking is when to use the base type of the type mark. The bounds of a range are often of type UNIVERSAL INTEGER, FIXED or FLOAT and in these instances the base type must be used. This is not correct when variables of other types are used to define a range.

#### SUNDRY FUNCTIONS OF THE TOKEN MODULE.

Sundry functions of the TOKEN MODULE include: procedures giving order to enumeration types; checking valid use of procedure and entry names; storing the bounds for arrays and static expressions; evaluating expressions for storage in DIANA nodes; finding the type of number declarations; counting array bounds; evaluating and storing the expression after the DIGITS or DELTA in accuracy constraints; assigning the type and range of loop parameters; keeping note of whether an array is dynamic and has enough bounds specified for it.

## 6. CHAPTER 6. THE DIANA MODULE.

DIANA is the intermediate code used to create the interface of the UCT-ADA front end with the rest of the ADA compiler or ADA tools that may be designed. The code produced is in the form of an abstract syntax tree where each of the nodes of the tree have predefined attributes indicating certain semantic or syntactic properties.

This tree is represented by a linked list where each element of the list is a node in the tree. Each node has the same basic structure. The attributes of the nodes have various meanings depending on their node type.

APPENDIX IIa and IIb illustrate an example of some ADA fragments and their equivalent DIANA abstract syntax tree.

## 6.1. THE UCT IMPLEMENTATION OF DIANA.

Each node is represented by a record in PLUS (the language used for this module) and uses fifteen words of storage. The record is subdivided in the following manner:

- 1) Box-type. This is a nine bit integer used by the implementor to group alternate nodes with the same function in a structure. It is of no use to anyone else.
- 2) Box-kind. This nine bit integer is the unique number that identifies each node as a specific node type. Interpretation of this is vital to any user of the front end.
- 3) Reference. There are twelve of these eighteen bit integer fields present in each record. They represent pointers to other nodes in the tree. The first has the particular function of giving the node its unique sequence number for reference by other

nodes.

4) Value. Four one word value fields are provided to store any numeric values. They are defined to be of type real and any user of the DIANA tree can determine if the value stored was an integer by looking at the other attributes of the node.

5) Boolean. This is a one bit integer field that occurs six times. The first four occurrences are linked to the value field. Every time a value is stored the corresponding boolean is set. This is necessary as the value field is initialised to zero. If the boolean is set the value was actually stored, even if it is zero. The remaining bits are for attributes that are stored as true or false.

6) Rep. Four one word logical fields store the representation of each token in the form it is input. If the token contains more than 16 characters, nodes will be linked together to store the remaining characters.

#### 6.1.1. INTERPRETATION OF EACH NODE.

The box-type is of no use to anyone but the implementor so may be ignored. Appendix III gives a list of nodes (according to the DIANA MANUAL [ Goos 1981 ] ) with their corresponding unique box-kind number, and the structure of each node showing what references map to which attributes of each DIANA node.

Certain attributes have been catered for but not implemented yet so will not be mentioned. They are the attributes that contain certain semantic information, information that is machine dependent (i.e. cd-attributes in the manual) and any attributes that require a boolean result. The comment attribute and source position attribute have not been implemented but have been

provided for. The comment attribute is not mentioned whereas the source position is in APPENDIX III.

The references with particular meanings are:

One: This is the sequence number of the node and is generated by the program. Each new node gets a unique sequence number. These start from eleven and sequence numbers one to ten are used for implementation characteristics e.g. predefined type integer, float, string etc.

Eleven: If the token being stored has more than sixteen characters, special nodes are created to hold the subsequent groups of sixteen characters. These are linked through reference eleven of successive nodes. The number of nodes created will be sufficient to store the complete token.

Twelve: In ADA there may be lists of things. e.g. IDENTIFIER\_LIST. These correspond to lists in DIANA. Each list has a particular list header. The header node links to the first element of the list through reference two, and this element links to the next through reference twelve. All remaining elements of the list are linked through reference twelve.

APPENDIX IIc illustrates the DIANA representation produced by the UCT-front end for the example given in APPENDIX IIa and IIb. DIANA nodes are produced as the ADA syntax demands. The calls to the DIANA producing module are interspersed in the parser and static semantic checking routines. Once a syntax or semantic error is found, the production of DIANA stops.

The DIANA MODULE has an array of size fifty. (This is set by a constant 'NO-OF-BOXES'). As each node is produced, it is stored in this array and kept there until it has been completed. Completed means that all possible attributes (structural and semantic) have been assigned a value, and any other node that

should be linked to it, either has been linked to it or its sequence number has been stored somewhere for reference. Once a node is complete it may be written out. Nodes are generally written out a structure at a time. E.g. The simple assignment statement

```
A := 3;
```

may generate the following nodes:

- 1) ASSIGN node.
- 2) USED-NAME-NODE node.
- 3) NUMERIC-LITERAL node.

Nodes two and three above may have been complex structures. E.g.

```
A.B := 3*4+5 ;
```

If this had been the statement, it would have been pointless to write out the box for A and 3, before completing the box for B and \*4+5. The ASSIGN node will not be completed until the next statement is reached because it must be linked to the next statement. Writing out a "structure" means that when the next statement (or termination point) is reached, all nodes making up the previous statement or structure are written out.

Once a node has been written out, the array element that held that node is cleared to be re-used. The array thus functions as a circularly linked list.

Appendix IIc illustrates that the order in which the nodes of the DIANA tree are generated is haphazard when compared with the logical position of the nodes in the tree. This often implies that a tree is produced without there being a need for it. No other nodes reference it, and it just floats.

For example, as each part of an expression is parsed, a tree is constructed and the sequence number of the root is stored. A new root for the whole expression is created as the tree grows. This

new root must replace the old root. Once the complete expression has been parsed, that root must be linked as the beginning of the expression.

This is accomplished through a stack which keeps track of the root of the tree being dealt with. As a new root is formed, all references to the old root must be changed to point to the new root. All references to the expression must be to the root of the whole tree, which may only be known once the tree is completed.

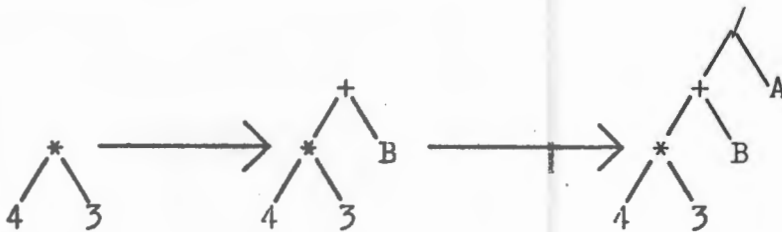


Figure 6-1.  $(4*3+B)/A$

## 6.2. STACKS AND STRUCTURES USED IN THE DIANA MODULE.

One of the design principles of DIANA is "regularity" so structures produced are often very similar. ADA allows nesting of structures within each other. This makes scope a very important factor in the production of DIANA. Several different stacks have been used to overcome the problem of scope.

Stacks play a vital role in the DIANA module. They are used to prevent linking of structures across nesting levels, block levels and to aid linkage when complex structures are dealt with.

Every list in DIANA has a list header. Once the header has been linked to the first item of its list, there is no need to keep the header node in memory, but the position of the header node is required so that the entire list may be written out once the structure is complete. There is a stack which keeps the box-kind

of the headers that have been written out and their sequence numbers. Once their corresponding structures have been written out, they are removed from the stack.

Blocks may be nested within one another. The attributes and structures of the outermost block will be completed only when the inner blocks are complete. Linkage of nodes is not always in the backward direction, i.e. some attributes cannot be completed until further nodes have been generated. The number of these further nodes may be random, resulting in the necessity to search through the as yet incomplete nodes, find the required node and link it to the node just generated. There may be several incomplete nodes of the same kind, each belonging to a different block in the list. As the list is circular and nodes are not written out until complete, there is no guarantee that the first node found is the correct one.

The nesting of blocks is proper nesting. This was used to create a block-stack which holds the sequence number of the first node generated in each block. Every time a node is searched for (once the correct node kind is found) a check is done to ensure that the sequence number of the node is greater than or equal to the first sequence number of the block. Once a block is exited, all remaining nodes of that block are written out using the sequence number stored in the block-stack for that block. The stack top is then moved down by one.

The same idea is used with a statement stack. Each sequence of statements is linked together, and these may also be nested. Consider the following example. The line numbers are for reference.

```
      BEGIN
1     C := 0 ;
2     B := 0 ;
3     A := 3 ;
4     WHILE A < 10 LOOP
5         A := A + 1;
6         IF (A * 2) > 10 THEN
7             B := B + 1;
8             C := C + 2;
9         ELSE
10            B := B + 2;
11           C := C + 1;
12        END IF;
13    END LOOP ;
14    IF C > B THEN
15        A := C;
16    ELSE
17        A := B ;
18    END IF;
19    D := 10;
20    END ;
```

Figure 6-2.

In this example there are several sequences of statements. The following line numbers form the different sequences.

- 1) 1, 2, 3, 4, 14, 19
- 2) 5, 6
- 3) 7, 8
- 4) 10, 11
- 5) 15
- 6) 17

The most complex stack used is the one called "seq-no-store". ADA as a language is ambiguous. Many of the ambiguities were overcome by restructuring the syntax, or by using look-ahead. Consequently, in many productions a node is generated before the actual structure being parsed has been defined. The structure header is therefore only generated in the middle of the structure and must still be linked to its components, and other nodes must be linked to it. The production of DIANA often assumes that the next node to be generated will be the header node of the following structure.

Wherever this problem occurs in the ADA syntax, use is made of the "seq-no-store" stack. Productions are nested so a stack is ideal. Each time a production is entered and a node is generated before its header, the sequence number of this node is stored on the stack, and the stack top incremented. When the header is finally generated, the sequence number on the stack indicates the first component of that structure, and it also indicates the sequence number which has been referenced instead of the header. This means that all references to this sequence number can be changed to refer to the current header. All references on the stack are also changed so that if a new header should occur higher in the nesting structure, the stack will have a record of the existing header. This mechanism is used in the building of expression trees and name trees.

In expression, at each level of the parse tree, the stack is incremented and the sequence number of the next node to be generated is stored. As the parse tree backtracks, the stack is decremented level by level. The same is done with name.

In the following example of name, the kind of name is only determined at the '(' or '.'. E.g. A(F.W(I))

The DIANA nodes are generated in the following way:

SEQUENCE NOS. ON STACK	ADA TOKEN	NODE SEQ-NO	TYPE OF NODE	TREE DEVELOPMENT
1	A	1	USED-NAME-ID	A
1	(	2	INDEXED	INDEXED ├── A
1 3	F	3	USED-NAME-ID	INDEXED ├── A └── F
1 4	.	4	SELECTED	INDEXED ├── A └── SELECTED └── F
1 4	W	5	USED-NAME-ID	INDEXED ├── A └── SELECTED ├── F └── W
1 6	(	6	INDEXED	INDEXED ├── A └── INDEXED └── SELECTED ├── F └── W
1 6	I	7	USED-NAME-ID	INDEXED ├── A └── INDEXED └── SELECTED ├── F └── W └── I
1	)			
	)			

Figure 6-3.

Stacks are therefore as important to the DIANA module as the token module is to the front end. All the problems created by the scope rules, nesting, and by ambiguities, have been solved using stacks.

## 6.3. STRUCTURE OF THE MODULE.

The DIANA module is made up of two major procedures with several others supporting them. Each procedure is one large case statement, where each case may be called from outside the module or within the module.

Procedure "genrate-diana-boxes" is the case statement which produces the nodes. It also calls the routines which manipulate the circularly linked list to determine where to put the newly created node.

Each unique box-kind is not produced by one particular constituent case. Several cases may produce the same box-kind but they may have different box-types, or acquire their attributes at different times, or be a particular box-kind that is produced in two different contexts.

The main function of this routine is to produce and complete the DIANA nodes concurrently as far as possible. Certain attributes can only be completed once several other nodes have been produced. Therefore there must be a further procedure that enables this procedure to search backwards through the nodes, find a particular node and complete the required attribute. The procedure that does this searches for a particular box-kind and then checks to see that the node found is in the correct block.

The second major procedure is called procedure "polishing". Its purpose is to do any "tidying up" that is needed. The cases making up this case statement are called from particular places where special processing is required. When structures need to be written out, a call will be made to this routine which will find the required header node, see what its sequence number is and call the writing out routine. Certain attributes can be completed at

various points in the ADA syntax. At these points, this routine will be called and the relevant information completed.

All the stack incrementing and decrementing and setting and resetting of flags is done through this case statement. Semantic values are received and transferred to the relevant nodes through procedure "polishing".

If any attributes require alteration, this routine finds the relevant nodes and alters the fields. This is mainly used where the ADA structure cannot be identified immediately, and a basic structure is created. Once the actual structure has been identified, the basic structure is changed.

Consider the ADA syntax for the IDENTIFIER\_LIST of a NUMBER\_DECLARATION, OBJECT\_DECLARATION, EXCEPTION\_DECLARATION, COMPONENT\_LIST and DISCRIMINANT\_LIST. It is identical in all of the above examples and consequently they all generate the IDENTIFIER\_LIST structure. DIANA differentiates between them all by generating NUMBER-ID, VAR-ID, EXCEPTION-ID, COMPONENT-ID and DISCRIMINANT-ID nodes. All these nodes have the same basic structure but are different box-kinds. At the end of each list, when it can be determined in the ADA syntax with which list is being dealt, procedure "polishing" is called and the relevant case finds the list header and changes all components of that list to their correct unique box-kind.

#### 6.4. SEMANTIC INFORMATION.

An important feature of the DIANA tree is the semantic information it holds. The UCT-front end does the static semantic checking using the ADA symbol table and the results are passed to the DIANA structure (through procedure "polishing"). As this information

comes from the ADA symbol table, the symbol table must be visible at all times.

The UCT-front end uses the symbol table to store certain DIANA attributes that may be referenced by other nodes further down the tree once the original nodes have been output.

The attributes held in the symbol table used by the DIANA module are sequence numbers which point to various characteristics of a definition. They correspond to the following semantic attributes:

- 1) SM-DEFN. The sequence number where a particular identifier is declared.
- 2) SM-TYPE-STRUCTURE. Type structure refers to the node giving the structure of the type e.g. record, integer, array. Each time a new type is defined, sm-type-struct is set. For derived types or access types, it is the structure of the first type definition.
- 3) SM-BASE-TYPE. Contrary to type structure, this attribute is that of the last type declaration in a sequence of declarations.
- 4) SM-CONSTRAINT. This attribute is that of the most recent constraint imposed on the particular identifier.

The base type of the range node of a range or range constraint is also obtained from information held in the ADA symbol table.

#### 6.5. OUTPUT OF THE DIANA TREE.

Each structure is written out as it is completed. This output is to a sequential disc file which must then be sorted with the key being each node sequence number for further use, or be used as a random access file. Nodes belonging to a particular structure will not necessarily be together on the file. In the formation of a list, as the nodes are linked, they are written out to avoid overwriting a link that already exists. Due to nesting, although

the body of a block may be written out, the header would only be written out much later when the next block at the same level was encountered.

Each record is output to a file as a string of characters. The record length is 120 characters. As PLUS has no random access facilities, a sequential file was produced.

The internal representation for DIANA used by the UCT-front end is in accordance with the specifications of the DIANA manual [Goos 1981]. The representation used for the DIANA nodes is similar to the structure of the nodes as specified in the manual and it would therefore be possible to write a driver routine which could translate the representation used into any desired form.

## 7. CHAPTER 7. WORKED EXAMPLES USING THE UCT-ADA FRONT END.

## 7.1. EXAMPLE 1.

This simple example illustrates the three areas covered by the UCT-front end.

Figure 7-1 is the program run with syntax errors. The error message indicates what the syntax error is and where error recovery begins.

Figure 7-2 is the program run with deliberate semantic errors. The first error is the simplest form of type incompatibility. The second is also a type mismatch but is not an obvious error. I-1.0 is legal, float - float. The error occurs because the function FACT returns an integer and therefore I\*FACT( ) is illegal i.e. float \* integer.

Figure 7-3 is the correct program compiled producing the DIANA given in figure 7-4.

Figure 7-1.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/20/83 12:46:09
  1  1  0  --TEST2
  2  1  0  FUNCTION FACT(I: INTEGER) RETURN INTEGER IS
  3  2  0  BEGIN
  4  2  0      IF I = 1
  5  2  0          RETURN 1 ;
***ERROR*** 182 'THEN' EXPECTED , RESERVED-WORD RETURN AT
LINE 5 COL 5 SKIPS AND RECOVERS TO RETURN AT LINE 5 COL 5
  6  2  0      ELSE
  7  2  0          RETURN I*FACT(I-1) ;
  8  2  0      END IF
  9  2  0  END FACT ;
***ERROR*** 202 ';' EXPECTED , RESERVED-WORD END AT
LINE 9 COL 1 SKIPS AND RECOVERS TO END AT LINE 9 COL 1
***** END PROCESSING 2 ERRORS

```

Figure 7-2.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/20/83 12:49:37
  1  1  0  --TEST2
  2  1  0  FUNCTION FACT(I: FLOAT) RETURN INTEGER IS
  3  2  0  BEGIN
  4  2  0      IF I = 1 THEN
***ERROR*** 504 OPERANDS OF DIFFERENT TYPES.
  5  2  0          RETURN 1 ;
  6  2  0      ELSE
  7  2  0          RETURN I*FACT(I-1.0) ;
***ERROR*** 510 OPERANDS OF ILLEGAL TYPES FOR THIS OPERATION .
  8  2  0      END IF ;
  9  2  0  END FACT ;
***** END PROCESSING 2 ERRORS

```



0	29	37	38	39	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	134	39	40	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	109	47	49	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	126	40	42	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	157	41	18	0	0	0	0	0	0	0	0
	45.00000		.00000		00I						
0	62	42	43	44	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	138	43	0	0	0	0	0	0	0	0	0
	0.00000		.00000		00*						
0	109	44	41	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	62	45	46	47	0	4	0	0	0	0	0
	0.00000		.00000		00						
0	157	46	14	0	0	0	0	0	0	0	0
	0.00000		.00000		00FACT						
0	157	48	18	0	0	0	0	0	0	0	0
	52.00000		.00000		00I						
0	62	49	50	51	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	138	50	0	0	0	0	0	0	0	0	0
	0.00000		.00000		00-						
0	109	51	48	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	100	52	4	0	0	0	0	0	0	0	0
	0.10000+001		.00000		101						
0	71	27	28	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	157	29	18	0	0	0	0	0	0	0	0
	33.00000		.00000		00I						
0	62	30	31	32	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	138	31	0	0	0	0	0	0	0	0	0
	0.00000		.00000		00=						
0	109	32	29	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	100	33	4	0	0	0	0	0	0	0	0
	0.10000+001		.00000		101						
0	166	38	0	0	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	19	25	0	26	0	0	0	0	0	0	0
	0.00000		.00000		00						
0	27	12	0	0	24	0	0	0	0	0	0
	0.00000		.00000		00						

## 7.2. EXAMPLE 2.

This example introduces interesting features of ADA, that of an incomplete type declaration, and the access type declaration. Default initial values are specified for the record LEADER and a function is defined and called within the program.

Figure 7-5 illustrates syntactic and semantic errors. Notice that because of the error in the definition of TREF, all variables of type TREF are not recognised as record variables. As a result of the method used to store records, the record components can only be accessed through the record variable name resulting in the generation of several unnecessary error messages.

Figure 7-6 is the syntactically correct program and APPENDIX IVa is the DIANA output by the UCT-front end.

Figure 7-5.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/21/83 22:12:43
  1  1  0  --EX91--
  2  1  0  PROCEDURE TSORT IS --TOPOLOGICAL SORT
  3  1  0  TYPE LEADER ;
  4  2  0  TYPE TRAILER ;
  5  2  0  TYPE LREF IS ACCESS LEADER
  6  2  0  TYPE TREF IS ACCESS ;
***ERROR*** 202 ';' EXPECTED , RESERVED-WORD TYPE AT
LINE 6 COL 3 SKIPS AND RECOVERS TO TYPE AT LINE 6 COL 3
***ERROR*** 260 <TYPE MARK> EXPECTED , SPECIAL ; AT
LINE 6 COL 24 SKIPS AND RECOVERS TO ; AT LINE 6 COL 24
  7  2  0  TYPE LEADER IS
  8  2  0  RECORD
  9  3  0  KEY      : INTEGER      := 0;
 10  3  0  COUNT   : INTEGER      := 0;
 11  3  0  TRAIL   : TREF         := NULL;
 12  3  0  NXT     : LREF         := NULL;
 13  3  0  END RECORD;
 14  2  0  TYPE TRAILER IS
 15  2  0  RECORD
 16  4  0  ID      : LREF ;
 17  4  0  NXT     : TREF ;
 18  4  0  END RECORD ;
 19  2  0  HEAD, TAIL, P, Q : LREF;
 20  2  0  T        : TREF;
 21  2  0  Z        : INTEGER ;
 22  2  0  X, Y     : INTEGER;
 23  2  0  FUNCTION L(W:INTEGER) RETURN LREF IS
 24  5  0  H : LREF ;
 25  5  0  BEGIN
 26  5  0  H := HEAD ;
 27  5  0  TAIL.KEY := H;
***ERROR*** 508 ASSIGNMENT ATTEMPTED BETWEEN OPERANDS OF DIFFERENT
TYPES .
 28  5  0  WHILE TAIL.KEY /= T LOOP
***ERROR*** 504 OPERANDS OF DIFFERENT TYPES.
 29  6  0  H := H.NXT;
 30  6  0  END LOOP;
 31  5  0  IF H = TAIL
 32  5  0  TAIL := NEW LEADER;
***ERROR*** 182 'THEN' EXPECTED , VARIABLE-ID TAIL AT
LINE 32 COL 8 SKIPS AND RECOVERS TO TAIL AT LINE 32 COL 8
 33  5  0  Z := Z + 1;
 34  5  0  H.COUNT := 0;
 35  5  0  H.TRAIL := NULL ;
 36  5  0  H.NXT := TAIL;
 37  5  0  END IF ;
 38  5  0  RETURN H;
 39  5  0  END L;
 40  2  0  BEGIN
 41  2  0  HEAD := NEW LEADER;
 42  2  0  TAIL := HEAD;
 43  2  0  Z := 0;
 44  2  0  GET(X);
 45  2  0  WHILE X /= 0 LOOP

```

```

46 7 0      GET(Y) ;
47 7 0      PUT(X); PUT(Y);
48 7 0      P := L(1.1);
***ERROR*** 534 FORMAL AND ACTUAL PARAMETERS MIS-MATCHED
49 7 0      Q := L(Y);
50 7 0      T := NEW TRAILER(ID => NULL, NXT = NULL);
***ERROR*** 510 OPERANDS OF ILLEGAL TYPES FOR THIS OPERATION .
51 7 0      T.ID := Q;
***ERROR*** 156 UNRECOGNIZABLE STATEMENT , IDENTIFIER ID AT
LINE 51 COL 7 SKIPS AND RECOVERS TO ID AT LINE 51 COL 7
***ERROR*** 528 UNDECLARED IDENTIFIER .
52 7 0      T.NXT := P.TRAIL;
***ERROR*** 156 UNRECOGNIZABLE STATEMENT , IDENTIFIER NXT AT
LINE 52 COL 7 SKIPS AND RECOVERS TO NXT AT LINE 52 COL 7
***ERROR*** 528 UNDECLARED IDENTIFIER .
53 7 0      P.TRAIL := T;
54 7 0      Q.COUNT := Q.COUNT +1;
55 7 0      GET(X);
56 7 0      END LOOP;
57 2 0      P := HEAD;
58 2 0      HEAD := NULL;
59 2 0      WHILE P /= TAIL LOOP
60 8 0      Q := P ;
61 8 0      P := P.NXT;
62 8 0      IF Q.COUNT = 0 THEN
63 8 0      Q.NXT := HEAD;
64 8 0      HEAD := Q;
65 8 0      END IF ;
66 8 0      END LOOP ;
67 2 0      Q := HEAD ;
68 2 0      OUTER : WHILE Q /= NULL LOOP
69 9 0      PUT(Q.KEY);
70 9 0      Z := Z - 1;
71 9 0      T := Q.TRAIL;
72 9 0      Q := Q.NXT;
73 9 0      INNER : WHILE T /= NULL LOOP
74 10 0     P := T.ID;
***ERROR*** 156 UNRECOGNIZABLE STATEMENT , IDENTIFIER ID AT
LINE 74 COL 21 SKIPS AND RECOVERS TO ID AT LINE 74 COL 21
***ERROR*** 532 INVALID ENTRY CALL OR PROCEDURE CALL.
75 10 0     P.COUNT := P.COUNT -1;
76 10 0     IF P.COUNT = 0 THEN
77 10 0     P.NXT := Q ;
78 10 0     Q := P;
79 10 0     END IF;
80 10 0     T := T.NXT;
***ERROR*** 156 UNRECOGNIZABLE STATEMENT , IDENTIFIER NXT AT
LINE 80 COL 21 SKIPS AND RECOVERS TO NXT AT LINE 80 COL 21
***ERROR*** 532 INVALID ENTRY CALL OR PROCEDURE CALL.
81 10 0     END LOOP INNER;
82 9 0      END LOOP OUTER;
83 2 0      IF Z = 0 THEN
84 2 0      PUT("THIS SET IS NOT PARTIALLY ORDERED.");
85 2 0      END IF ;
86 2 0      END TSORT;
***** END PROCESSING 15 ERRORS

```

Figure 7-6.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/21/83 22:24:16
 1 1 0 --EX91--
 2 1 0 PROCEDURE TSORT IS --TOPOLOGICAL SORT
 3 1 0 TYPE LEADER ;
 4 2 0 TYPE TRAILER ;
 5 2 0 TYPE LREF IS ACCESS LEADER ;
 6 2 0 TYPE TREF IS ACCESS TRAILER ;
 7 2 0 TYPE LEADER IS
 8 2 0 RECORD
 9 3 0 KEY : INTEGER := 0;
10 3 0 COUNT : INTEGER := 0;
11 3 0 TRAIL : TREF := NULL;
12 3 0 NXT : LREF := NULL;
13 3 0 END RECORD;
14 2 0 TYPE TRAILER IS
15 2 0 RECORD
16 4 0 ID : LREF ;
17 4 0 NXT : TREF ;
18 4 0 END RECORD ;
19 2 0 HEAD, TAIL, P, Q : LREF;
20 2 0 T : TREF;
21 2 0 Z : INTEGER ;
22 2 0 X, Y : INTEGER;
23 2 0 FUNCTION L(W:INTEGER) RETURN LREF IS
24 5 0 H : LREF ;
25 5 0 BEGIN
26 5 0 H := HEAD ;
27 5 0 TAIL.KEY := W;
28 5 0 WHILE TAIL.KEY /= W LOOP
29 6 0 H := H.NXT;
30 6 0 END LOOP;
31 5 0 IF H = TAIL
32 5 0 THEN
33 5 0 TAIL := NEW LEADER;
34 5 0 Z := Z + 1;
35 5 0 H.COUNT := 0;
36 5 0 H.TRAIL := NULL ;
37 5 0 H.NXT := TAIL;
38 5 0 END IF ;
39 5 0 RETURN H;
40 5 0 END L;
41 2 0 BEGIN
42 2 0 HEAD := NEW LEADER;
43 2 0 TAIL := HEAD;
44 2 0 Z := 0;
45 2 0 GET(X);
46 2 0 WHILE X /= 0 LOOP
47 7 0 GET(Y) ;
48 7 0 PUT(X); PUT(Y);
49 7 0 P := L(X);
50 7 0 Q := L(Y);
51 7 0 T := NEW TRAILER(ID => NULL, NXT = NULL);
52 7 0 T.ID := Q;
53 7 0 T.NXT := P.TRAIL;
54 7 0 P.TRAIL := T;

```

```

55 7 0      Q.COUNT := Q.COUNT +1;
56 7 0      GET(X);
57 7 0      END LOOP;
58 2 0      P := HEAD;
59 2 0      HEAD := NULL;
60 2 0      WHILE P /= TAIL LOOP
61 8 0          Q := P ;
62 8 0          P := P.NXT;
63 8 0          IF Q.COUNT = 0 THEN
64 8 0              Q.NXT := HEAD;
65 8 0              HEAD := Q;
66 8 0          END IF ;
67 8 0      END LOOP ;
68 2 0      Q := HEAD ;
69 2 0      OUTER : WHILE Q /= NULL LOOP
70 9 0          PUT(Q.KEY);
71 9 0          Z := Z - 1;
72 9 0          T := Q.TRAIL;
73 9 0          Q := Q.NXT;
74 9 0      INNER : WHILE T /= NULL LOOP
75 10 0          P := T.ID;
76 10 0          P.COUNT := P.COUNT -1;
77 10 0          IF P.COUNT = 0 THEN
78 10 0              P.NXT := Q ;
79 10 0              Q := P;
80 10 0          END IF;
81 10 0          T := T.NXT;
82 10 0      END LOOP INNER;
83 9 0      END LOOP OUTER;
84 2 0      IF Z = 0 THEN
85 2 0          PUT("THIS SET IS NOT PARTIALLY ORDERED.");
86 2 0      END IF ;
87 2 0      END TSORT;
***** END PROCESSING 0 ERRORS

```

## 7.3. EXAMPLE 3.

This example illustrates that the syntax checker validates the syntax for most of ADA and not just the "PASCAL" subset. An interesting feature of this program is the use of an aggregate to initialise the values of the array of type WORK-WEEK.

Figure 7-7.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/21/83 22:44:57
  1  1  0  PACKAGE WORK DATA IS
  2  1  0      TYPE DAY IS (MON,TUES,WED,THU,FRI,SAT,SUN);
  3  1  0      TYPE TIME IS DELTA 0.01 RANGE 0.0 .. 24.0 ;
  4  1  0      TYPE WORK WEEK IS ARRAY(MON .. SUN) OF TIME ;
  5  1  0      NORMAL HOURS : CONSTANT WORK WEEK := (MON .. FRI
              => 8.0 , SAT .. SUN => 0.0)
  6  1  0  END ;
  7  1  0  FUNCTION OVERTIME(ACTUAL: WORK_WEEK)
              RETURN FLOAT IS
  8  2  0  OVER : FLOAT := 0.0;
  9  2  0  I : INTEGER ;
 10  2  0  BEGIN
 11  2  0      FOR I IN DAY LOOP
 12  3  0          OVER := OVER + FLOAT(ACTUAL(I))
              - FLOAT(NORMAL_HOURS(I))
 13  3  0      END LOOP;
 14  2  0      RETURN OVER ;
 15  2  0  END OVERTIME ;
 16  1  0  FUNCTION SALARY(RATE: FLOAT ;
              HOURS : WORK_WEEK)RETURN FLOAT IS
 17  4  0  PAY, EXCESS : FLOAT ;
 18  4  0  BEGIN
 19  4  0      EXCESS := OVERTIME(HOURS);
 20  4  0      IF EXCESS < 0.0 THEN
 21  4  0          PAY := (40.0 + EXCESS) * RATE;
 22  4  0      ELSE
 23  4  0          PAY := 40.0*RATE + 2.0*EXCESS*RATE;
 24  4  0      END IF;
 25  4  0      RETURN PAY;
 26  4  0  END SALARY;
***** END PROCESSING 0 ERRORS

```

## 7.4. EXAMPLE 4.

This example again includes features not covered by the semantic and code generation phases of this thesis. It was used to check that loop parameters could be used in expressions with the loop parameter having the correct default type.

Figure 7-8.

```
@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/21/83 22:58:19
 1  1  0  PACKAGE ARRAYTEST IS PRIVATE END ;
 2  1  0  PACKAGE BODY ARRAYTEST IS
 3  2  0  LOOPARRAY : ARRAY ( 1 .. 10 ) OF INTEGER ;
 4  2  0  J : INTEGER ;
 5  2  0  BEGIN
 6  2  0  FOR I IN 1 .. 10 LOOP
 7  3  0  LOOPARRAY (I) := I ;
 8  3  0  END LOOP ;
 9  2  0  FOR J IN 3 .. 12 LOOP
10  4  0  LOOPARRAY ( J - 3 ) := LOOPARRAY ( J -3)
      + LOOPARRAY(J)
11  4  0  END LOOP ;
12  2  0  END ;
***** END PROCESSING 0 ERRORS
```

## 7.5. EXAMPLE 5.

Figure 7-9 illustrates other errors that may occur.

Figure 7-10 is the corrected program and the DIANA produced may be found in APPENDIX IVb.

Figure 7-9.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/21/83 23:21:29
  1  1  0  --EX90--
  2  1  0  PROCEDURE INSERT IS --LINEAR LIST INSERTION.
  3  1  0  TYPE NODE;
  4  2  0  TYPE REF IS ACCESS NODE ;
  5  2  0  TYPE NODE IS
  6  2  0  RECORD
  7  3  0  KEY : INTEGER ;
  8  3  0  COUNT : INTEGER ;
  9  3  0  NXT : REF ;
 10  3  0  END RECORD ;
 11  2  0  K : INTEGER ;
 12  2  0  ROOT : REF := NEW NODE(0,0,NULL) ;
 13  2  0  PROCEDURE SEARCH(X : INTEGER ; BASE :
        IN OUT REF) IS
 14  4  0  W : REF ;
 15  4  0  B : BOOLEAN ;
 16  4  0  BEGIN
 17  4  0  W := BASE ;
 18  4  0  B := TRUE ;
 19  4  0  WHILE W LOOP
***ERROR*** 514 CONDITION MUST BE A BOOLEAN EXPRESSION.
 20  5  0  IF W.KEY = X
 21  5  0  THEN B := FALSE ;
 22  5  0  ELSE W := W.NXT ;
 23  5  0  END IF ;
 24  5  0  END LOOP ;
 25  4  0  IF B THEN -- NEW ENTRY
 26  4  0  W := BASE ;
 27  4  0  BASE := NEW NODE(0,0,NULL) ;
 28  4  0  BASE.KEY := X ;
 29  4  0  BASE.COUNT := 1.1 ;
***ERROR*** 508 ASSIGNMENT ATTEMPTED BETWEEN OPERANDS OF DIFFERENT
TYPES .
 30  4  0  ELSE
 31  4  0  W.COUNT := W.COUNT + 1 ;
 32  4  0  END IF ;
 33  4  0  END SEARCH ;
 34  2  0  PROCEDURE PRINTLIST(V:REF) IS
 35  6  0  BEGIN
 36  6  0  WHILE V /= NULL LOOP
 37  7  0  PUT(V.KEY);
 38  7  0  PUT(V.COUNT);
 39  7  0  V := V.NXT ;
 40  7  0  END LOOP ;

```

## EXAMPLE 5.

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```
41 6 0 END PRINTLIST;
42 2 0 BEGIN -- MAIN PROGRAM
43 2 0 GET(K);
44 2 0 WHILE K /= 0
45 8 0 SEARCH(K,ROOT);
***ERROR*** 150 'LOOP' EXPECTED , PROCEDURE-ID SEARCH AT
LINE 45 COL 5 SKIPS AND RECOVERS TO SEARCH AT LINE 45 COL 5
46 8 0 GET(K) ;
47 8 0 END LOOP ;
48 2 0 PRINTLIST(ROOT) ;
49 2 0 END INSERT ;
***** END PROCESSING 3 ERRORS
```

Figure 7-10.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/21/83 23:04:33
 1 1 0 --EX90--
 2 1 0 PROCEDURE INSERT IS --LINEAR LIST INSERTION.
 3 1 0 TYPE NODE;
 4 2 0 TYPE REF IS ACCESS NODE ;
 5 2 0 TYPE NODE IS
 6 2 0 RECORD
 7 3 0 KEY : INTEGER ;
 8 3 0 COUNT : INTEGER ;
 9 3 0 NXT : REF ;
10 3 0 END RECORD ;
11 2 0 K : INTEGER ;
12 2 0 ROOT : REF := NEW NODE(O,O,NULL) ;
13 2 0 PROCEDURE SEARCH(X : INTEGER ; BASE :
    IN OUT REF) IS
14 4 0 W : REF ;
15 4 0 B : BOOLEAN ;
16 4 0 BEGIN
17 4 0 W := BASE ;
18 4 0 B := TRUE ;
19 4 0 WHILE W /= NULL AND B LOOP
20 5 0 IF W.KEY = X
21 5 0 THEN B := FALSE ;
22 5 0 ELSE W := W.NXT ;
23 5 0 END IF ;
24 5 0 END LOOP ;
25 4 0 IF B THEN -- NEW ENTRY
26 4 0 W := BASE ;
27 4 0 BASE := NEW NODE(O,O,NULL) ;
28 4 0 BASE.KEY := X ;
29 4 0 BASE.COUNT := 1 ;
30 4 0 ELSE
31 4 0 W.COUNT := W.COUNT + 1 ;
32 4 0 END IF ;
33 4 0 END SEARCH ;
34 2 0 PROCEDURE PRINTLIST(V:REF) IS
35 6 0 BEGIN
36 6 0 WHILE V /= NULL LOOP
37 7 0 PUT(V.KEY);
38 7 0 PUT(V.COUNT);
39 7 0 V := V.NXT ;
40 7 0 END LOOP ;
41 6 0 END PRINTLIST;
42 2 0 BEGIN -- MAIN PROGRAM
43 2 0 GET(K);
44 2 0 WHILE K /= 0 LOOP
45 8 0 SEARCH(K,ROOT);
46 8 0 GET(K) ;
47 8 0 END LOOP ;
48 2 0 PRINTLIST(ROOT) ;
49 2 0 END INSERT ;
***** END PROCESSING 0 ERRORS

```

## 7.6. EXAMPLE 6.

EXAMPLE 6 is not a sensible example and was produced to illustrate the DIANA produced for sequences of statements. This may be found in APPENDIX IVc.

Figure 7-11.

```

@RELOCATES.FINAL1,SL
ADA SCANNER VERSION U.C.T 1983 03/21/83 23:26:19
 1 1 0  PROCEDURE V IS
 2 1 0  TYPE INTEGER IS RANGE 1 .. 50 ;
 3 2 0  TYPE COLOUR IS (RED, ORANGE, YELLOW, GREEN) ;
 4 2 0  ROBOT : COLOUR ;
 5 2 0  VAR2 , VAR1 : INTEGER ;
 6 2 0  BEGIN
 7 2 0  <<LAB1>> VAR1 := 2 ;
 8 2 0  <<LAB2>> <<LAB3>> VAR2 := VAR1 ;
 9 2 0  VAR2 := 10 ;
10 2 0  NULL ;
11 2 0  GOTO LAB2 ;
12 2 0  GOTO LAB1 ;
13 2 0  GOTO LAB3 ;
14 2 0  <<LAB4>> NAMED1 : LOOP
15 3 0  CASE VAR1 IS
16 3 0  WHEN 1\2 =>
17 3 0  DECLARE
18 4 0  VAR4 : FLOAT ;
19 4 0  BEGIN
20 4 0  VAR4 := 4.5 * 3.3 ;
21 4 0  VAR4 := VAR4 ** 2 ;
22 4 0  END ;
23 3 0  WHEN 3 => LOOP
24 5 0  VAR1 := VAR1 + 1 ;
25 5 0  EXIT WHEN VAR1 = 10 ;
26 5 0  END LOOP ;
27 3 0  WHEN OTHERS => NULL ;
28 3 0  END CASE ;
29 3 0  VAR1 := VAR2 ;
30 3 0  VAR2 := 20 ;
31 3 0  EXIT NAMED1 ;
32 3 0  END LOOP NAMED1 ;
33 2 0  NAMED2 : FOR I IN COLOUR LOOP
34 6 0  VAR1 := VAR1 + 1 ;
35 6 0  VAR2 := VAR2 + VAR1 ;
36 6 0  EXIT NAMED2 WHEN VAR1 = 3 ;
37 6 0  END LOOP ;
38 2 0  NAMED3 : DECLARE
39 7 0  VAR3 : FLOAT ;
40 7 0  BEGIN
41 7 0  VAR3 := 3.3 ;
42 7 0  IF VAR1 = VAR2
43 7 0  THEN
44 7 0  VAR3 := 6.6 ;
45 7 0  END IF ;

```

```

46 7 0      END ;
47 2 0      WHILE VAR1 < 10 LOOP
48 8 0      IF VAR2 = 10
49 8 0      THEN
50 8 0      VAR1 := 5 ;
51 8 0      ELSIF VAR2 = 9
52 8 0      THEN
53 8 0      VAR1 := 7 ;
54 8 0      ELSIF VAR2 = 7
55 8 0      THEN
56 8 0      VAR1 := 5 ;
57 8 0      ELSE
58 8 0      VAR1 := 3 ;
59 8 0      END IF ;
60 8 0      EXIT ;
61 8 0      END LOOP ;
62 2 0      FOR J IN REVERSE GREEN .. RED LOOP
63 9 0      ROBOT := J ;
64 9 0      RETURN ;
65 9 0      RETURN VAR1 ;
66 9 0      NULL ;
67 9 0      END LOOP ;
68 2 0      ROBOT := GREEN ;
69 2 0      END ;
***** END PROCESSING 0 ERRORS

```

## 8. CHAPTER 8. CONCLUSION.

The purpose of this thesis was the production of a front end processor for a subset of ADA. This goal has been achieved for the "PASCAL" subset of ADA with several other structures included but not explicitly mentioned. The processor runs on a Univac 1100 system and was produced by using a compiler generator package for the syntax checker, PLUS routines for the semantic analysis and code generation. It produces an internal form of DIANA consistent with the DIANA specification.

This output is being used by several honours students at the University of Cape Town who are designing programming environment tools for ADA.

Experience has been gained in using the compiler generator package GSA 1100, and many of the problems encountered by the implementors at the University of Cape Town have been subsequently reported in the literature [ Bonet ] [ Harrison 1982 ] [ Winkler 1981 ]. It appears that the solutions used by other implementors coincide favourably with those used in the UCT-front end processor.

In undertaking this project greater insight has been achieved into the ADA programming language, DIANA, and the area of program environments. The project provides a basis for further research.



=> CHARACTER\_STRING \

```

$[DIGIT]$ ( # '_' [ $[DIGIT]$ \ <* 30 => ILLEGAL ] \ $[DIGIT]$ )
[ [ '#' \ # ':' $ '#' ] [ $[ $[DIGIT]$ $[BASEARITH]$ ]$
\ <* 35 => ILLEGAL ]
( # '_' [ $[ $[DIGIT]$ $[BASEARITH]$ ]$ \ <* 30 => ILLEGAL ] \
$[ $[DIGIT]$ $[BASEARITH]$ ]$ ) ...
?[ '.' '..' : => BASED_NUMBER ]?
( '.' [ $[ $[DIGIT]$ $[BASEARITH]$ ]$ \ <* 35 => ILLEGAL ]
( # '_' [ $[ $[DIGIT]$ $[BASEARITH]$ ]$
\ <* 30 => ILLEGAL ] \ $[ $[DIGIT]$ $[BASEARITH]$ ]$ ) ... )
[ [ '#' \ # ':' $ '#' ] \ <* 36 => ILLEGAL ]
( ! $[ 'e' 'E' ]$ ( $[SIGN]$ ) [ $[DIGIT]$ \ <* 35 => ILLEGAL ]
( # '_' [ $[DIGIT]$ \ <* 30 => ILLEGAL ] \ $[DIGIT]$ ) ... )
=> BASED_NUMBER \
?[ '.' '..' : => DECIMAL_NUMBER ]?
( '.' [ $[DIGIT]$ \ <* 35 => ILLEGAL ]
( # '_' [ $[DIGIT]$ \ <* 30 => ILLEGAL ] \ $[DIGIT]$ ) ... )
( ! $[ 'E' 'e' ]$ ( $[SIGN]$ ) [ $[DIGIT]$ \ <* 35 => ILLEGAL ]
( # '_' [ $[DIGIT]$ \ <* 30 => ILLEGAL ] \ $[DIGIT]$ ) ... ) ]
=> DECIMAL_NUMBER \

```

```

'-' [ '-' # ( $[COMMENTCHARS]$ ) ...
[ EOL -> TOKEN \ ** <* 50 -> TOKEN ] \ => SPECIAL ] ] .

```

```

$EJECT
$[SIGN]$ ::= $[ '+' '-' ]$ .
$[LETTER]$ ::= $[ 'a' .. 'z' 'A' .. 'Z' ]$ .
$[DIGIT]$ ::= $[ '0' .. '9' ]$ .
$[SPECIALS]$ ::= $[ '&' ')' '+' ',' ';' '|' ]$ .
$[STRINGCHARS]$ ::= $[ NUL .. '!' '#' .. DEL ]$ .
$[ILLEGALS]$ ::= $[ NUL .. SP '#' .. '%' '?' .. '@' '[' .. '\ '
'{' '}' .. EOL EOC .. BAD ]$ .
$[COMMENTCHARS]$ ::= $[ NUL .. DEL ]$ .
$[BASEARITH]$ ::= $[ 'A' 'B' 'C' 'D' 'E' 'F' ]$ .
$[CHARACTERLIT]$ ::= $[ '!' .. '~' ]$ .

```

## 10. APPENDIX II. EXAMPLE OF DIANA PRODUCTION.

## 10.1. APPENDIX IIa. ADA FRAGMENTS.

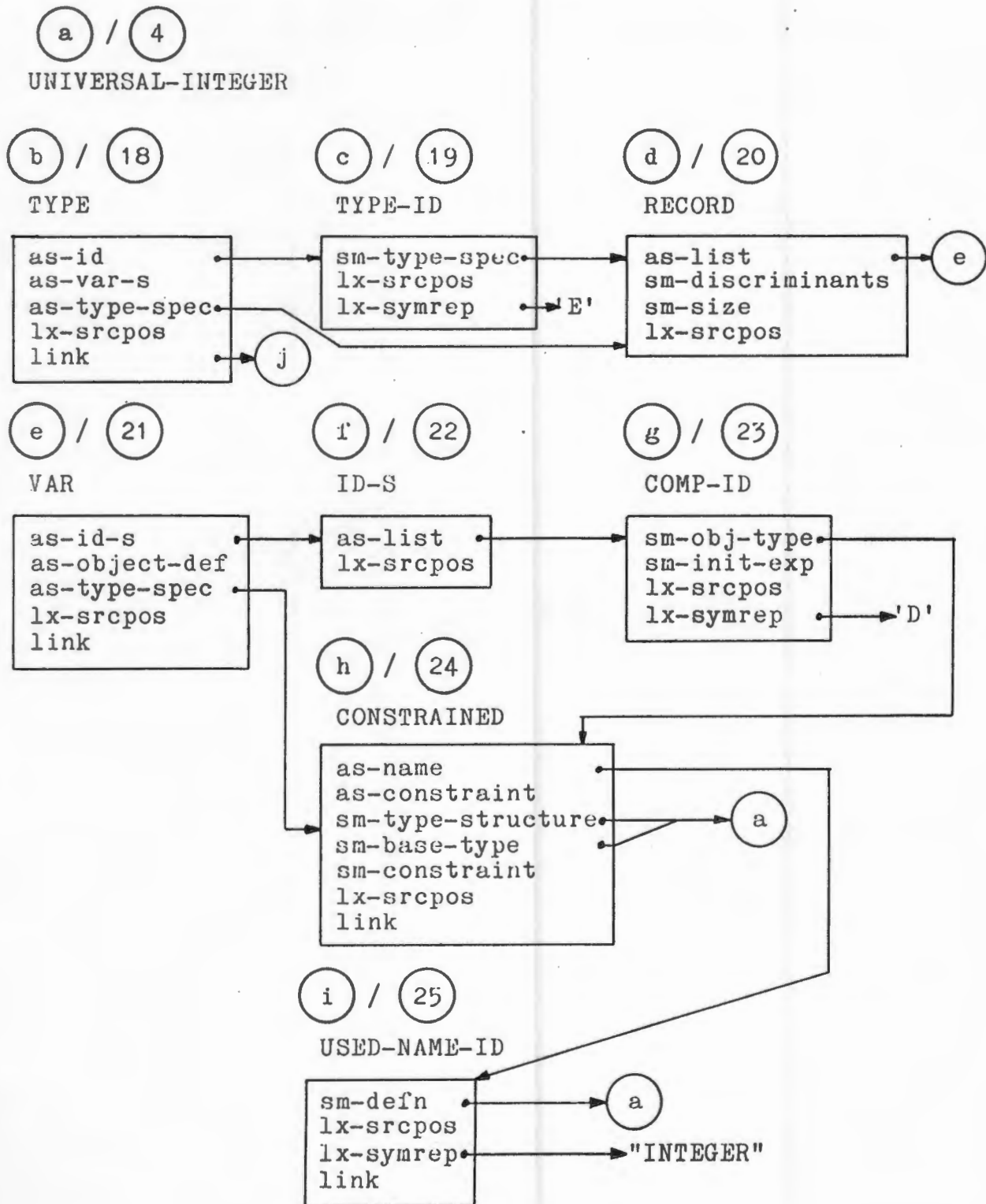
```
.  
. .  
TYPE E IS RECORD  
    D : INTEGER ;  
END RECORD ;  
C : E ;  
A,B : INTEGER ;  
F : ARRAY ( 1 .. 10 ) OF INTEGER ;  
BEGIN  
    A := 3 ;  
    B := 0 ;  
    WHILE B < 10 LOOP  
        F(B) := (4*3+B)/A ;  
        B := B+A ;  
    END LOOP ;  
    C.D := B ;  
END  
. . .
```

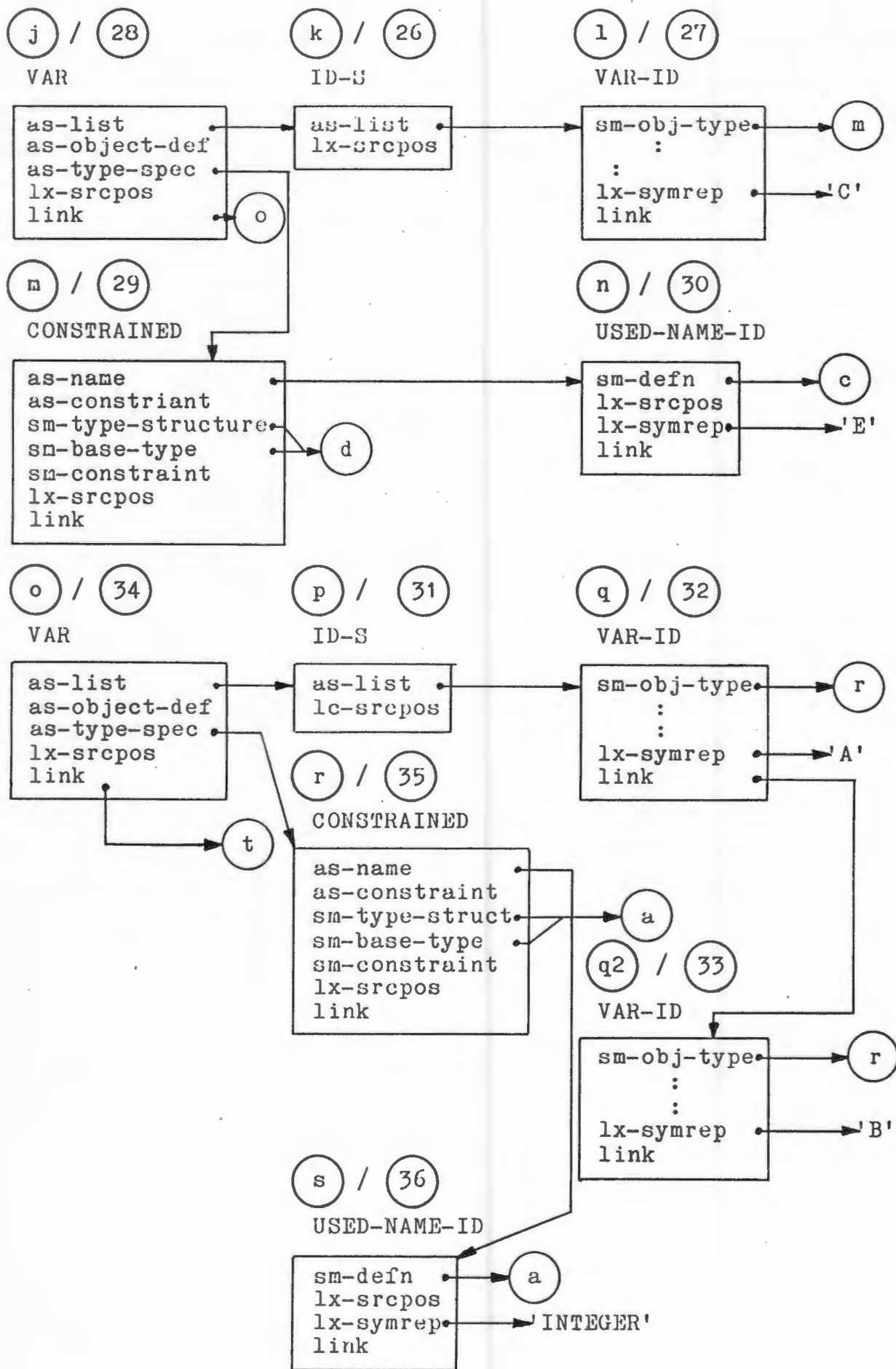
10.2. APPENDIX IIB. DIANA ABSTRACT SYNTAX TREE FRAGMENTS.

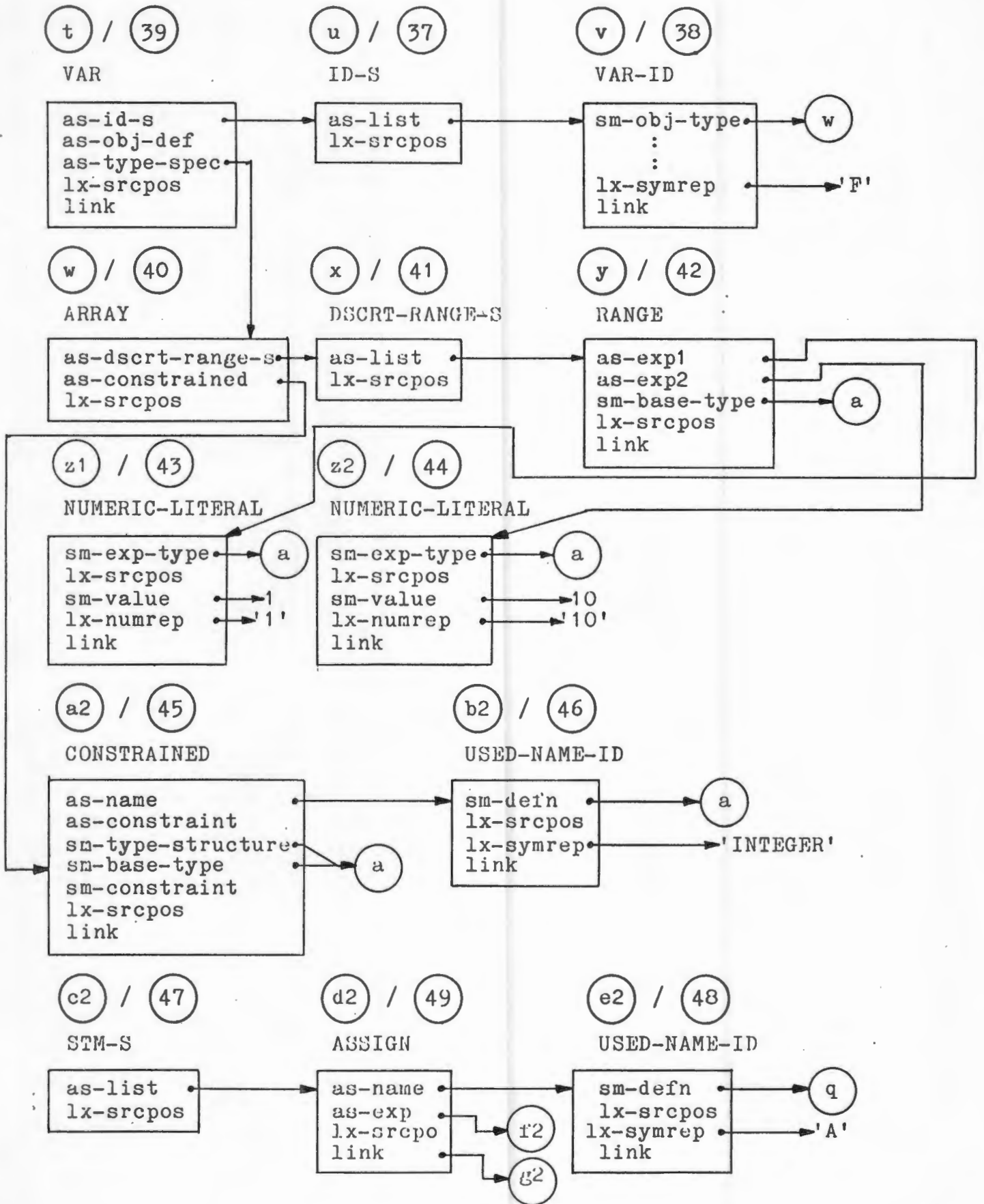
Key (a) : connector

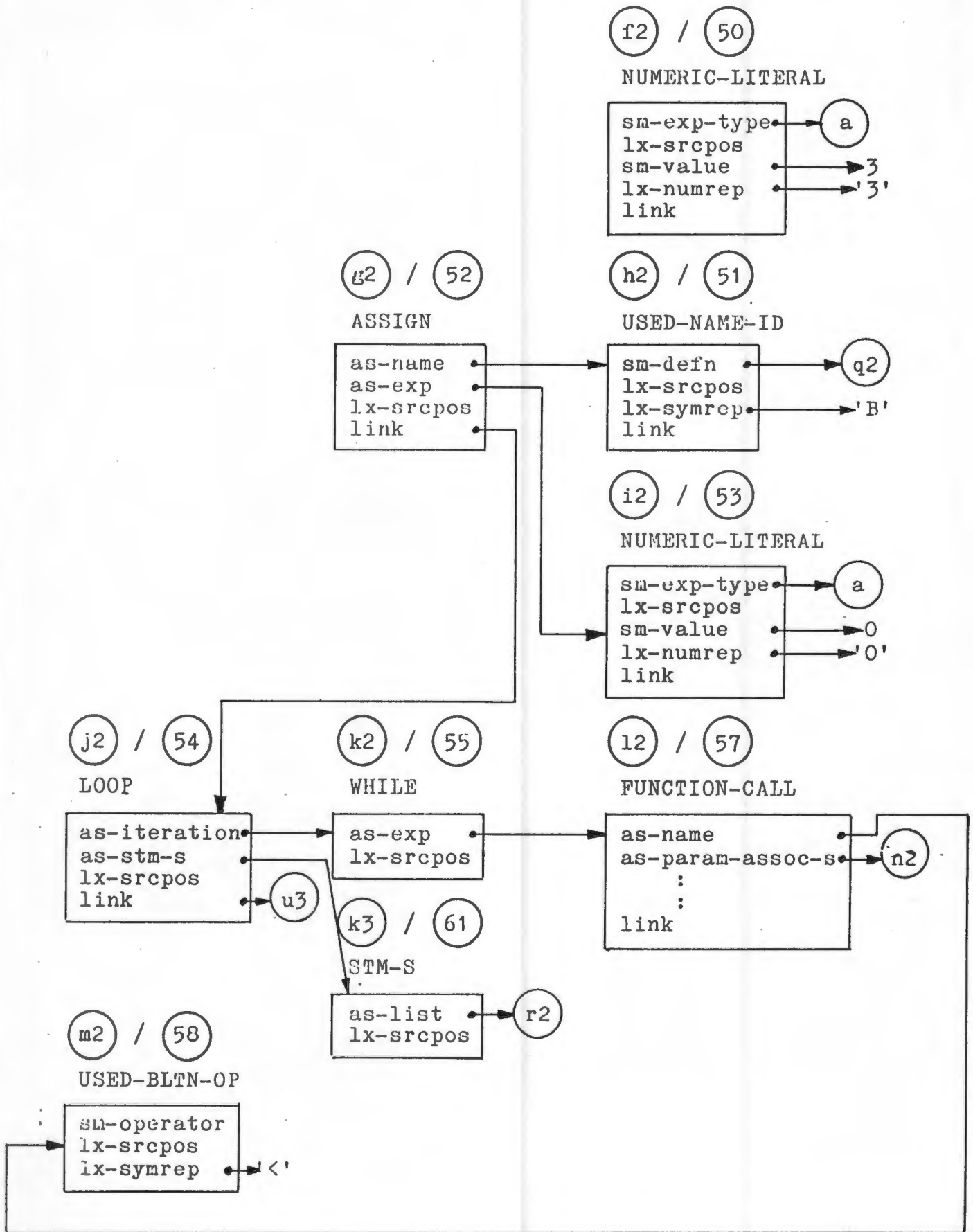
Note : Alphabetic node numbers - hand generated sequence nos.

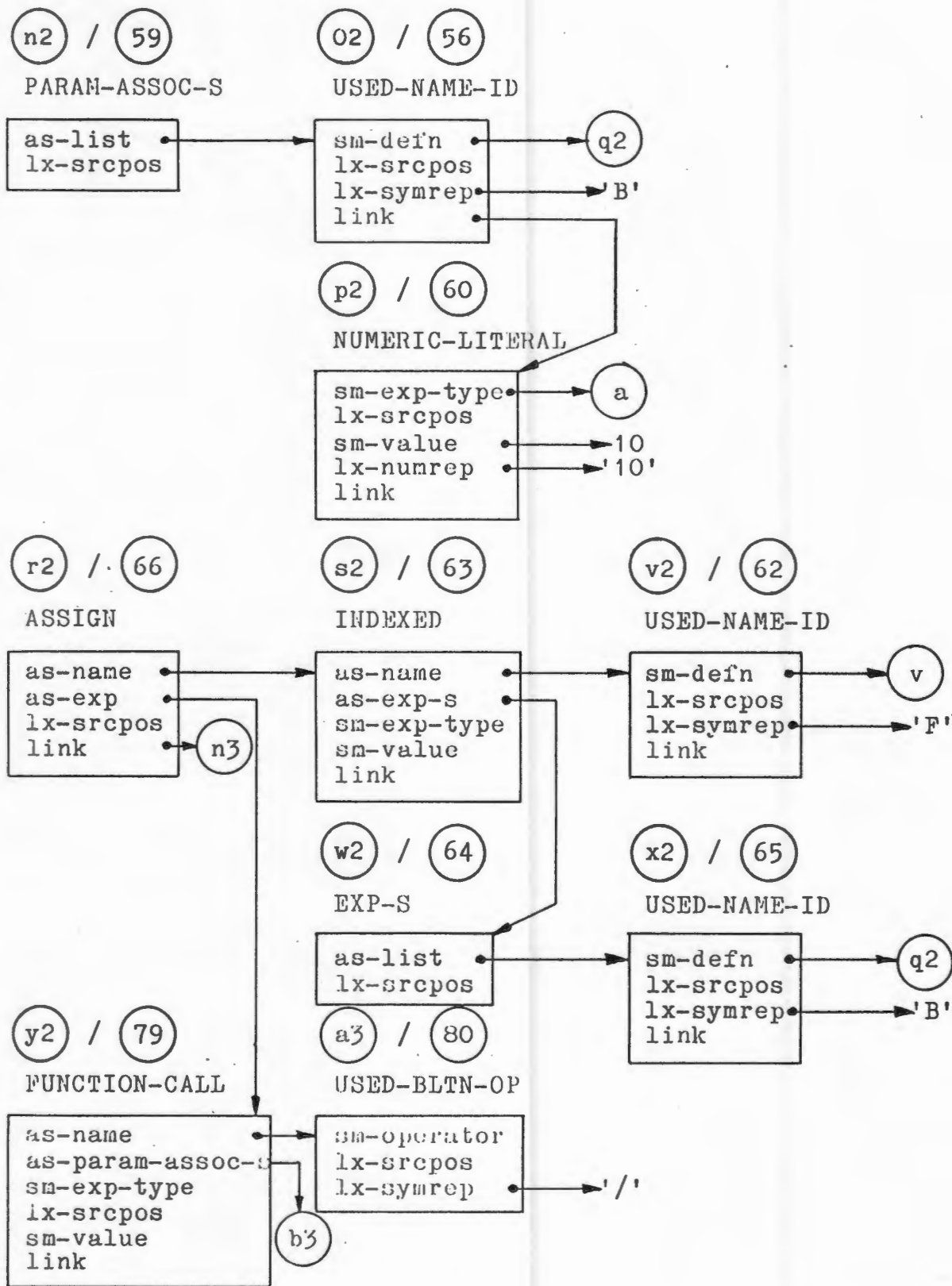
Numeric node numbers - program generated sequence nos.

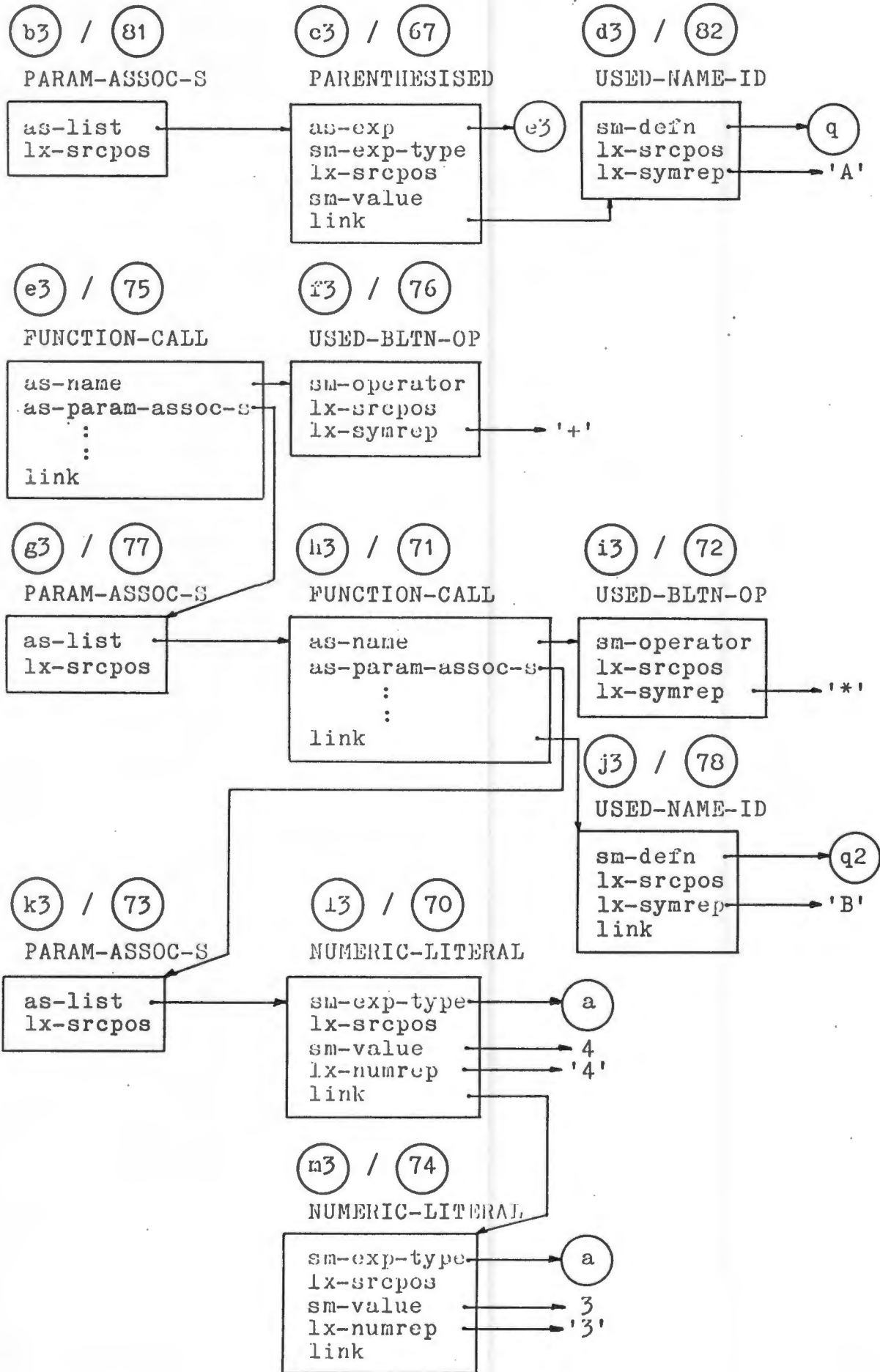


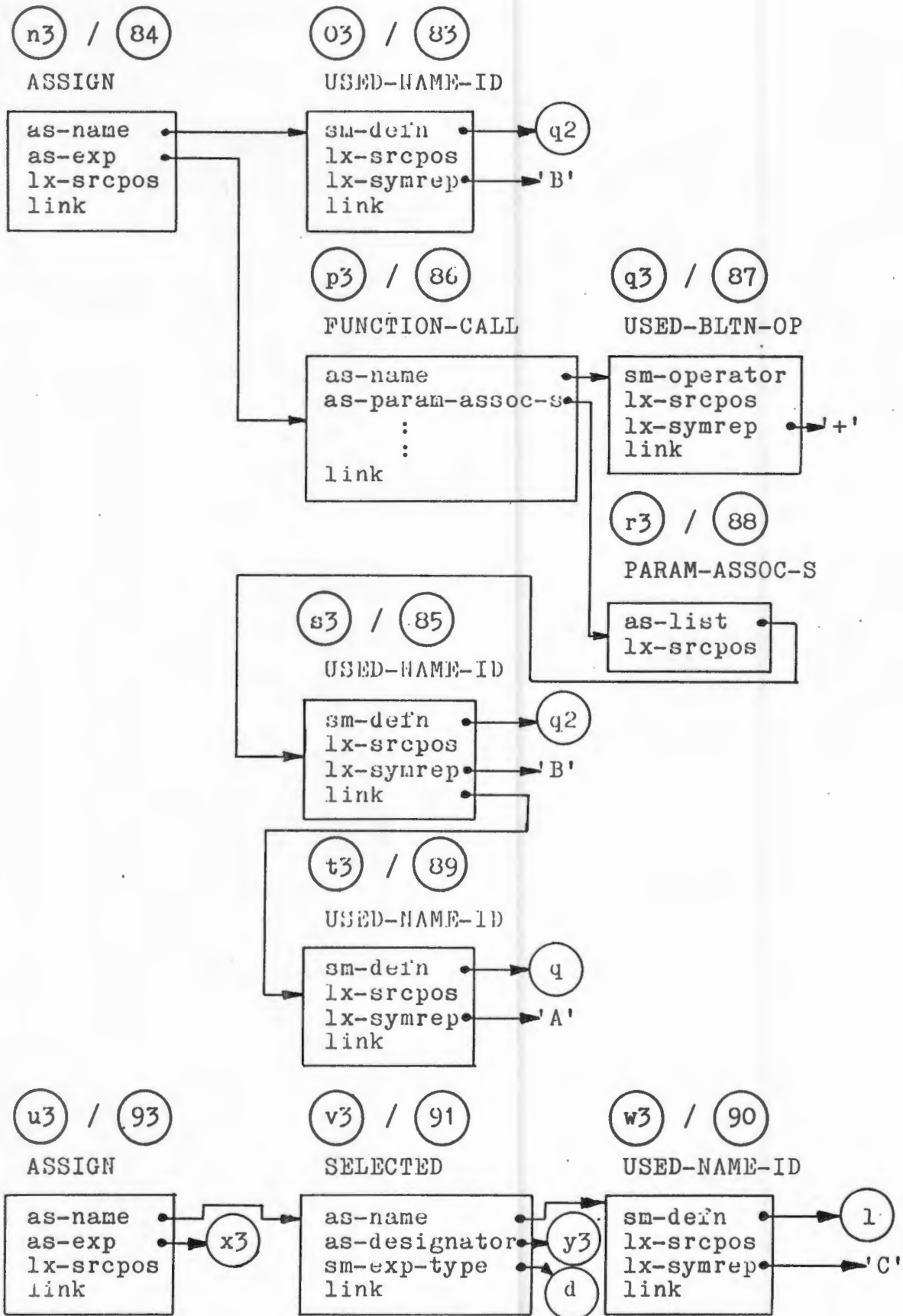


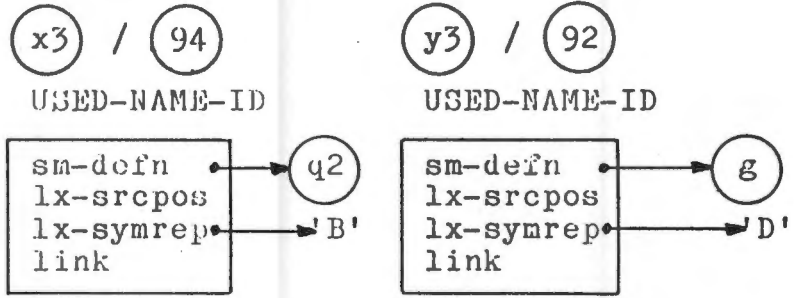
















APPENDIX II. EXAMPLE OF DIANA PRODUCTION.

KND:	111	100	62	138	109	100	62	138	109	157
SEQ:	67	70	71	72	73	74	75	76	77	78
REF2:	75	4	72		70	4	76		71	33
REF3:			73				77			
REF4:										
REF5:										
REF6:										
REF11:										
REF12:	82	74	78							
VAL1:		4				3	3			
REP1:		4		*		3		+		B
REP2:										

KND:	62	138	109	157	157	13	157	62	138	109
SEQ:	79	80	81	82	83	84	85	86	87	88
REF2:	80		67	32	33	83	33	87		85
REF3:	81					86		88		
REF4:										
REF5:										
REF6:										
REF11:										
REF12:							89			
VAL1:		/								
REP1:		/		A	B		B		+	
REP2:										

KND:	157	157	131	157	13	157
SEQ:	89	90	91	92	93	94
REF2:	32	27	90	23	91	33
REF3:			92		94	
REF4:			20			
REF5:						
REF6:						
REF11:						
REF12:						
VAL1:						
REP1:	A	C		D		B
REP2:						

## 11. APPENDIX III. STRUCTURE OF DIANA NODES.

## Note:

- 1) reference(1) is always the sequence number of the node, i.e. node generation number. It is not mentioned in this appendix.
- 2) The attribute "link" does not appear in the manual, but is used to link the nodes forming a list.

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
ABORT	1		not implemented
ACCEPT	2		not implemented
ACCESS	3	as-constrained sm-size sm-storage-size lx-srcpos cd-impl-size  cd-alignment  sm-controlled	reference(2) reference(3) reference(4) reference(5) value(1) boolean(1) value(2) boolean(2) boolean(5)
ADDRESS	4		not implemented
AGGREGATE	5	as-list sm-exp-type sm-constraint lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(12) value(1) boolean(1)
ALL	6	as-name sm-exp-type lx-srcpos link sm-value	reference(2) reference(4) reference(5) reference(12) value(1) boolean(1)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
ALLOCATOR	7	as-name as-access-constraint sm-exp-type lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(12) value(1) boolean(1)
ALTERNATIVE	8	as-choice-s as-stm-s lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
ALTERNATIVE-S	9	as-list lx-srcpos	reference(2) reference(3)
AND-THEN	10	lx-srcpos	reference(2)
ARGUMENT-ID	11		not implemented
ARRAY	12	as-dscrt-range-s as-constrained sm-size lx-srcpos cd-impl-size  cd-alignment  sm-packing	reference(2) reference(3) reference(4) reference(5) value(1) boolean(1) value(2) boolean(2) boolean(5)
ASSIGN	13	as-name as-exp lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
ASSOC	14	as-id as-actual lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
ATTR-ID	15		not implemented
ATTRIBUTE	16		not implemented
ATTRIBUTE-CALL	17		not implemented

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
BINARY	18	as-exp1 as-binary-op as-exp2 sm-exp-type lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(6) reference(12) value(1) boolean(1)
BLOCK	19	as-item-s as-stm-s as-alternative-s lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)
BOX	20		not implemented
CASE	21	as-exp as-alternative-s lx-srcpos	reference(2) reference(3) reference(4)
CHOICE-S	22	as-list-s lx-srcpos	reference(2) reference(3)
CODE	23	as-exp lx-srcpos	reference(2) reference(3)
COMP-ID	24	sm-obj-type sm-init-exp lx-srcpos link cd-position  cd-first-bit  cd-last-bit  lx-symrep	reference(2) reference(3) reference(4) reference(12) value(1) boolean(1) value(2) boolean(2) value(3) boolean(3) rep(1) - rep(4)
COMP-REP	25		not implemented
COMP-REP-S	26		not implemented
COMP-UNIT	27	as-pragma-s as-context as-unit-body lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
COMPILATION	28	as-list lx-srcpos	reference(2) reference(3)
COND-CLAUSE	29	as-exp-void as-stm-s lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
COND-ENTRY	30		not implemented
CONST-ID	31	sm-obj-type sm-obj-def sm-address lx-srcpos link lx-symrep	reference(2) reference(3) reference(4) reference(5) reference(12) rep(1) - rep(4)
CONSTANT	32	as-id-s as-object-def as-type-spec lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)
CONSTRAINED	33	as-name as-constraint sm-type-struct sm-base-type sm-constraint lx-srcpos link cd-impl-size  cd-alignment	reference(2) reference(3) reference(4) reference(5) reference(6) reference(7) reference(12) value(1) boolean(1) value(2) boolean(2)
CONTEXT	34	as-list lx-srcpos	reference(2) reference(3)
CONVERSION	35	as-name as-exp sm-exp-type lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(12) value(1) boolean(1)
DECL-REP-S	36		not implemented
DECL-S	37	as-list lx-srcpos	reference(2) reference(3)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
DEF-CHAR	38	sm-obj-type lx-srcpos sm-pos  sm-rep  lx-symrep	reference(2) reference(3) value(1) boolean(1) value(2) boolean(2) rep(1) - rep(4)
DEF-OP	39	sm-spec sm-body sm-location lx-srcpos lx-symrep	reference(2) reference(3) reference(4) reference(5) rep(1) - rep(4)
DELAY	40		not implemented
DERIVED	41	as-constrained sm-size sm-storage-size lx-srcpos cd-impl-size  cd-alignment  sm-actual-delta  sm-packing sm-controlled	reference(2) reference(3) reference(4) reference(5) value(1) boolean(1) value(2) boolean(2) value(3) boolean(3) boolean(5) boolean(6)
DISCRIMANT-AGGREGATE	42	as-list lx-srcpos	reference(2) reference(3)
DSCRMT-ID	43	sm-obj-type sm-init-exp lx-srcpos cd-position  cd-first-bit  cd-last-bit  lx-symrep	reference(2) reference(3) reference(4) value(1) boolean(1) value(2) boolean(2) value(3) boolean(3) rep(1) - rep(4)
DSCRMT-RANGE-S	44	as-list lx-srcpos	reference(2) reference(3)
ENTRY	45		not implemented
ENTRY-CALL	46		not implemented

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
ENTRY-ID	47	sm-spec sm-address lx-srcpos lx-symrep	reference(2) reference(3) reference(4) rep(1) - rep(4)
ENUM-ID	48	sm-obj-type lx-srcpos sm-pos  sm-rep  lx-symrep	reference(2) reference(3) value(1) boolean(1) value(2) boolean(2) rep(1) - rep(4)
ENUM-LITERAL-S49		as-list sm-size lx-srcpos cd-impl-size  cd-alignment	reference(2) reference(3) reference(4) value(1) boolean(1) value(2) boolean(2)
EXCEPTION	50		not implemented
EXCEPTION-ID	51	sm-exception-def lx-srcpos lx-symrep	reference(2) reference(3) rep(1) - rep(4)
EXIT	52	as-name-void as-exp-void sm-stm lx-srcpos	reference(2) reference(3) reference(4) reference(5)
EXP-S	53	as-list lx-srcpos	reference(2) reference(3)
FIXED	54	as-exp as-range-void sm-size lx-srcpos cd-impl-size  cd-alignment  sm-actual-delta  sm-bits	reference(2) reference(3) reference(4) reference(5) value(1) boolean(1) value(2) boolean(2) value(3) boolean(3) value(4) boolean(4)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
FLOAT	55	as-exp as-range-void sm-type-struct sm-size lx-srcpos cd-impl-size  cd-alignment	reference(2) reference(3) reference(4) reference(5) reference(6) value(1) boolean(1) value(2) boolean(2)
FOR	56	as-id as-dscrt-range lx-srcpos	reference(2) reference(3) reference(4)
FORMAL-DSCRT	57		not implemented
FORMAL-FIXED	58		not implemented
FORMAL-FLOAT	59		not implemented
FORMAL-INTEGER	60		not implemented
FUNCTION	61	as-param-s as-constrained-void lx-srcpos	reference(2) reference(3) reference(4)
FUNCTION-CALL	62	as-name as-param-assoc-s sm-exp-type lx-srcpos sm-value	reference(2) reference(3) reference(5) reference(6) value(1) boolean(1)
FUNCTION-ID	63	sm-spec sm-body sm-location lx-srcpos lx-symrep	reference(2) reference(3) reference(4) reference(5) rep(1) - rep(4)
GENERIC	64		not implemented
GENERIC-ASSOC-S	65		not implemented
GENERIC-ID	66	sm-generic-param-s sm-spec sm-body lx-srcpos lx-symrep	reference(2) reference(3) reference(4) reference(5) rep(1) - rep(4)
GENERIC-OP	67		not implemented
GENERIC-PARAM-S	68		not implemented

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
GO TO	69	as-name lx-srcpos link	reference(2) reference(3) reference(12)
ID-S	70	as-list lx-srcpos	reference(2) reference(3)
IF	71	as-list lx-srcpos link	reference(2) reference(3) reference(12)
IN	72	as-id-s as-type-spec as-exp-void lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)
IN-ID	73	sm-obj-type sm-init-exp lx-srcpos link lx-symrep	reference(2) reference(3) reference(5) reference(12) rep(1) - rep(4)
IN-OP	74	lx-srcpos	reference(2)
IN-OUT	75	as-id-s as-type-spec as-exp-void lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)
IN-OUT-ID	76	sm-obj-type lx-srcpos link lx-symrep	reference(2) reference(3) reference(12) rep(1) - rep(4)
INDEX	77	as-name lx-srcpos link	reference(2) reference(3) reference(12)
INDEXED	78	as-name as-exp-s sm-exp-type link sm-value	reference(2) reference(3) reference(4) reference(12) value(1) boolean(1)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
INNER-RECORD	79	as-list lx-srcpos	reference(2) reference(3)
INSTANTIATION	80		not implemented
INTEGER	81	as-range sm-type-struct sm-size lx-srcpos cd-impl-size  cd-alignment	reference(2) reference(3) reference(4) reference(5) value(1) boolean(1) value(2) boolean(2)
ITEM-S	82	as-list lx-srcpos	reference(2) reference(3)
ITERATION-ID	83	sm-obj-type lx-srcpos lx-symrep	reference(2) reference(3) rep(1) - rep(4)
L-PRIVATE	84		not implemented
L-PRIVATE-TYPE-ID	85		not implemented
LABEL-ID	86	sm-stm lx-srcpos link lx-symrep	reference(2) reference(3) reference(12) rep(1) - rep(4)
LABELLED	87	as-id-s as-stm lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
LOOP	88	as-iteration as-stm-s lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
MEMBERSHIP	89	as-exp as-membership-op as-type-range sm-exp-type lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(6) reference(12) value(1) boolean(1)
NAME-S	90	as-list lx-srcpos	reference(2) reference(3)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
NAMED	91	as-choice-s as-exp lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
NAMED-STM	92	as-id as-stm lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
NO-DEFAULT	93		not implemented
HOT-IN	94	lx-srcpos	reference(2)
NULL-ACCESS	95	sm-exp-type lx-srcpos sm-value	reference(2) reference(3) value(1) boolean(1)
NULL-COMP	96	lx-srcpos	reference(2)
NULL-STM	97	lx-srcpos	reference(2)
NUMBER	98	as-id-s as-exp lx-srcpos link	reference(2) reference(3) reference(5) reference(12)
NUMBER-ID	99	sm-obj-type as-init-exp lx-srcpos link lx-symrep	reference(2) reference(3) reference(4) reference(12) rep(1) - rep(4)
NUMERIC-LITERAL	100	sm-exp-type lx-srcpos link sm-value  lx-numrep	reference(2) reference(2) reference(12) value(1) boolean(1) rep(1) - rep(4)
OR-ELSE	101	lx-srcpos	reference(2)
OTHERS	102	lx-srcpos link	reference(2) reference(12)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
OUT	103	as-id-s as-type-spec as-exp-void lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)
OUT-ID	104	sm-obj-type lx-srcpos link lx-symrep	reference(2) reference(3) reference(12) rep(1) - rep(4)
PACKAGE-BODY	105		not implemented
PACKAGE-DECL	106		not implemented
PACKAGE-ID	107	sm-spec sm-body sm-address lx-srcpos lx-symrep	reference(2) reference(3) reference(4) reference(5) rep(1) - rep(4)
PACKAGE-SPEC	108		not implemented
PARAM-ASSOC-S	109	as-list lx-srcpos	reference(2) reference(3)
PARAM-S	110	as-list lx-srcpos	reference(2) reference(3)
PARENTHESISED	111	as-exp sm-exp-type lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(12) value(1) boolean(1)
PRAGMA	112	as-id as-param-assoc-s lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
PRAGMA-ID	113	as-list lx-srcpos	reference(2) reference(3)
PRAGMA-S	114	as-list lx-srcpos	reference(2) reference(3)
PRIVATE	115		not implemented

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
PRIVATE-TYPE-ID	116		not implemented
PROC-ID	117	sm-spec sm-body sm-location lx-srcpos lx-synrep	reference(2) reference(3) reference(4) reference(5) rep(1) - rep(4)
PROCEDURE	118	as-param-s lx-srcpos	reference(2) reference(3)
PROCEDURE-CALL	119	as-name as-param-assoc-s lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
QUALIFIED	120	as-name as-exp sm-exp-type lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(12) value(1) boolean(1)
RAISE	121		not implemented
RANGE	122	as-exp2 as-exp2 sm-base-type lx-srcpos link	reference(1) reference(3) reference(4) reference(5) reference(12)
RECORD	123	as-list sm-discriminants sm-size lx-srcpos cd-impl-size  cd-alignment  sm-packing	reference(2) reference(3) reference(4) reference(5) value(1) boolean(1) value(2) boolean(2) boolean(5)
RECORD-REP	124		not implemented
RENAME	125		not implemented
RETURN	126	as-exp-void lx-srcpos link	reference(2) reference(3) reference(12)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
REVERSE	127	as-id as-dscrt-range lx-srcpos	reference(2) reference(3) reference(4)
SELECT	128		not implemented
SELECT-CLAUSE	129		not implemented
SELECT-CLAUSE-S	130		not implemented
SELECTED	131	as-name as-designator sm-exp-type lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(12) value(1) boolean(1)
SIMPLE-REP	132		not implemented
SLICE	133	as-name as-dscrt-range sm-exp-type sm-constraint lx-srcpos link sm-value	reference(2) reference(3) reference(4) reference(5) reference(6) reference(12) value(1) boolean(1)
STM-S	134	as-list lx-srcpos	reference(2) reference(3)
STRING-LITERAL	135	sm-exp-type sm-constraint lx-srcpos link sm-value  lx-symrep	reference(2) reference(3) reference(4) reference(12) value(1) boolean(1) rep(1) - rep(4)
STUB	136		not implemented
SUBPROGRAM-BODY	137	as-designator as-header as-block-stub lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
USED-BLTN-OP	138	sm-operator lx-srcpos lx-symrep	reference(2) reference(3) rep(1) - rep(4)
SUBPROGRAM-DECL	139		not implemented
SUBTYPE	140	as-id as-constrained lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
SUBTYPE-ID	141	sm-type-spec lx-srcpos lx-symrep	reference(2) reference(3) rep(1) - rep(4)
SUBUNIT	142		not implemented
TASK-BODY	143		not implemented
TASK-BODY-ID	144	sm-type-spec sm-body lx-srcpos lx-symrep	reference(2) reference(3) reference(4) rep(1) - rep(4)
TASK-DECL	145		not implemented
TASK-SPEC	146		not implemented
TERMINATE	147		not implemented
TIMED-ENTRY	148		not implemented
TYPE	149	as-id as-var-s as-type-spec lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)
TYPE-ID	150	sm-type-spec lx-srcpos lx-symrep	reference(2) reference(3) rep(1) - rep(4)
UNIVERSAL-FIXED	151	denoted by a 2 as the attribute value	
UNIVERSAL-INTEGGER	152	denoted by a 4 as the attribute value	
UNIVERSAL-REAL	153	denoted by a 3 as the attribute value	

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD.
USE	154	as-list lx-srcpos link	reference(2) reference(3) reference(12)
USED-BLT-ID	155	sm-operator lx-srcpos lx-symrep	reference(2) reference(3) rep(1) - rep(4)
USED-CHAR	156	sm-defn sm-exp-type lx-srcpos sm-value	reference(2) reference(3) reference(4) value(1) boolean(1)
USED-NAME-ID	157	sm-defn lx-srcpos link lx-symrep	reference(2) reference(3) reference(12) rep(1) - rep(4)
USED-OBJECT-ID	158	sm-defn sm-exp-type lx-srcpos link sm-value  lx-symrep	reference(2) reference(3) reference(4) reference(12) value(1) boolean(1) rep(1) - rep(4)
USED-OP	159	sm-defn link lx-symrep	reference(2) reference(12) rep(1) - rep(4)
VAR	160	as-id-s as-object-def as-type-spec lx-srcpos link	reference(2) reference(3) reference(4) reference(5) reference(12)
VAR-ID	161	as-obj-type sm-obj-def sm-address lx-srcpos link lx-symrep	reference(2) reference(3) reference(4) reference(5) reference(12) rep(1) - rep(4)
VAR-S	162	as-list lx-srcpos	reference(2) reference(3)
VARIANT	163	as-choice-s as-record lx-srcpos	reference(2) reference(3) reference(4)

NODE NAME	TYPE NUMBER	ATTRIBUTES	REFERENCE IN RECORD:
VARIANT-PART	164	as-name as-variant-s lx-srcpos link	reference(2) reference(3) reference(4) reference(12)
VARIANT-S	165	as-list lx-srcpos	reference(2) reference(3)
VOID	166		
WHILE	167	as-exp lx-srcpos	reference(2) reference(3)
WITH	168	as-list lx-srcpos link	reference(2) reference(3) reference(12)
CONTINUATION	512	link	reference(11)

12. APPENDIX IV. DIANA PRODUCED FOR EXAMPLES IN CHAPTER 7.

12.1. APPENDIX IVa.

0	28	11	12	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	118	13	0	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	117	14	13	16	0	0	0	0	0	0	0
0	0.00000		.00000		00TSORT						
0	137	15	14	13	16	0	0	0	0	0	0
0	0.00000		.00000		00						
0	82	17	18	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	149	18	19	0	0	0	0	0	0	0	0
0	20.00000		.00000		00						
0	150	19	0	0	0	0	0	0	0	0	0
0	0.00000		.00000		00LEADER						
0	149	20	21	0	0	0	0	0	0	0	0
0	22.00000		.00000		00						
0	150	21	0	0	0	0	0	0	0	0	0
0	0.00000		.00000		00TRAILER						
0	149	22	23	0	24	0	0	0	0	0	0
0	27.00000		.00000		00						
0	150	23	24	0	0	0	0	0	0	0	0
0	0.00000		.00000		00LREF						
0	3	24	25	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	33	25	26	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	157	26	19	0	0	0	0	0	0	0	0
0	0.00000		.00000		00LEADER						
0	149	27	28	0	29	0	0	0	0	0	0
0	32.00000		.00000		00						
0	150	28	29	0	0	0	0	0	0	0	0
0	0.00000		.00000		00TREF						
0	3	29	30	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	33	30	31	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	157	31	21	0	0	0	0	0	0	0	0
0	0.00000		.00000		00TRAILER						
0	123	34	35	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	70	36	37	0	0	0	0	0	0	0	0
0	0.00000		.00000		00						
0	24	37	38	40	0	0	0	0	0	0	0
0	0.00000		.00000		00KEY						
0	33	38	39	0	4	4	0	0	0	0	0
0	0.00000		.00000		00						
0	157	39	4	0	0	0	0	0	0	0	0
0	0.00000		.00000		00INTEGER						
0	100	40	4	0	0	0	0	0	0	0	0
0	0.00000		.00000		100						
0	160	35	36	40	38	0	0	0	0	0	0

0	41.00000	.00000	00							
0	70	42	43	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	24	43	44	46	0	0	0	0	0	0
0	0.00000	.00000	00COUNT							
0	33	44	45	0	4	4	0	0	0	0
0	0.00000	.00000	00							
0	157	45	4	0	0	0	0	0	0	0
0	0.00000	.00000	00INTEGER							
0	100	46	4	0	0	0	0	0	0	0
0	0.00000	.00000	100							
0	160	41	42	46	44	0	0	0	0	0
0	47.00000	.00000	00							
0	70	48	49	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	24	49	50	52	0	0	0	0	0	0
0	0.00000	.00000	00TRAIL							
0	33	50	51	0	0	29	0	0	0	0
0	0.00000	.00000	00							
0	157	51	28	0	0	0	0	0	0	0
0	0.00000	.00000	00TREF							
0	95	52	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	160	47	48	52	50	0	0	0	0	0
0	53.00000	.00000	00							
0	70	54	55	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	24	55	56	58	0	0	0	0	0	0
0	0.00000	.00000	00NXT							
0	33	56	57	0	0	24	0	0	0	0
0	0.00000	.00000	00							
0	157	57	23	0	0	0	0	0	0	0
0	0.00000	.00000	00LREF							
0	95	58	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	160	53	54	58	56	0	0	0	0	0
0	0.00000	.00000	00							
0	149	32	33	0	34	0	0	0	0	0
0	59.00000	.00000	00							
0	150	33	34	0	0	0	0	0	0	0
0	0.00000	.00000	00LEADER							
0	123	61	62	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	70	63	64	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	24	64	65	0	0	0	0	0	0	0
0	0.00000	.00000	00ID							
0	33	65	66	0	0	24	0	0	0	0
0	0.00000	.00000	00							
0	157	66	23	0	0	0	0	0	0	0
0	0.00000	.00000	00LREF							
0	160	62	63	0	65	0	0	0	0	0
0	67.00000	.00000	00							
0	70	68	69	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	24	69	70	0	0	0	0	0	0	0
0	0.00000	.00000	00NXT							
0	33	70	71	0	0	29	0	0	0	0
0	0.00000	.00000	00							

157	71	28	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOTREF						
160	67	68	0	70	0	0	0	0	0	0
0	0.00000	.00000		00						
149	59	60	0	61	0	0	0	0	0	0
0	77.00000	.00000		00						
150	60	61	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOTRAILER						
70	72	73	0	0	0	0	0	0	0	0
0	0.00000	.00000		00						
161	73	78	0	0	0	0	0	0	0	0
0	74.00000	.00000		OOHEAD						
161	74	78	0	0	0	0	0	0	0	0
0	75.00000	.00000		OOTAIL						
161	75	78	0	0	0	0	0	0	0	0
0	76.00000	.00000		00P						
161	76	78	0	0	0	0	0	0	0	0
0	0.00000	.00000		00Q						
160	77	72	0	78	0	0	0	0	0	0
0	82.00000	.00000		00						
33	78	79	0	0	24	0	0	0	0	0
0	0.00000	.00000		00						
157	79	23	0	0	0	0	0	0	0	0
0	0.00000	.00000		00LREF						
70	80	81	0	0	0	0	0	0	0	0
0	0.00000	.00000		00						
161	81	83	0	0	0	0	0	0	0	0
0	0.00000	.00000		00T						
160	82	80	0	83	0	0	0	0	0	0
0	87.00000	.00000		00						
33	83	84	0	0	29	0	0	0	0	0
0	0.00000	.00000		00						
157	84	28	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOTREF						
70	85	86	0	0	0	0	0	0	0	0
0	0.00000	.00000		00						
161	86	88	0	0	0	0	0	0	0	0
0	0.00000	.00000		00Z						
160	87	85	0	88	0	0	0	0	0	0
0	93.00000	.00000		00						
33	88	89	0	4	4	0	0	0	0	0
0	0.00000	.00000		00						
157	89	4	0	0	0	0	0	0	0	0
0	0.00000	.00000		00INTEGER						
70	90	91	0	0	0	0	0	0	0	0
0	0.00000	.00000		00						
161	91	94	0	0	0	0	0	0	0	0
0	92.00000	.00000		00X						
161	92	94	0	0	0	0	0	0	0	0
0	0.00000	.00000		00Y						
110	98	99	0	0	0	0	0	0	0	0
0	0.00000	.00000		00						
72	99	100	102	104	0	0	0	0	0	0
0	0.00000	.00000		00						
70	100	101	0	0	0	0	0	0	0	0
0	0.00000	.00000		00						
73	101	102	104	0	0	0	0	0	0	0
0	0.00000	.00000		00W						
33	102	103	0	4	4	0	0	0	0	0

0	0.00000	.00000	00							
157	103	4	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOINTEGER							
166	104	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
160	93	90	0	94	0	0	0	0	0	0
0	107.00000	.00000	00							
33	94	95	0	4	4	0	0	0	0	0
0	0.00000	.00000	00							
157	95	4	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOINTEGER							
61	96	98	105	0	0	0	0	0	0	0
0	0.00000	.00000	00							
63	97	96	108	0	0	0	0	0	0	0
0	0.00000	.00000	OOL							
33	105	106	0	0	24	0	0	0	0	0
0	0.00000	.00000	00							
157	106	23	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOLREF							
82	109	112	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
70	110	111	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
161	111	113	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOH							
134	115	117	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	116	111	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOH							
13	117	116	118	0	0	0	0	0	0	0
0	122.00000	.00000	00							
157	118	73	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOHEAD							
157	119	74	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOTAL							
131	120	119	121	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	121	37	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOKEY							
13	122	120	123	0	0	0	0	0	0	0
0	124.00000	.00000	00							
157	123	101	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOW							
134	133	135	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	134	111	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOH							
13	135	134	137	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	136	111	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOH							
131	137	136	138	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	138	55	0	0	0	0	0	0	0	0
0	0.00000	.00000	OONXT							
88	124	125	133	0	0	0	0	0	0	0
0	139.00000	.00000	00							
167	125	129	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							



0	0.00000	.00000	00							
157	169	111	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOH							
131	170	169	171	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	171	55	0	0	0	0	0	0	0	0
0	0.00000	.00000	OONXT							
13	172	170	173	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	173	74	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOTAIL							
71	139	140	0	0	0	0	0	0	0	0
0	174.00000	.00000	OO							
29	140	142	146	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	141	111	0	0	0	0	0	0	0	0
0	145.00000	.00000	OOH							
62	142	143	144	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
138	143	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO=							
109	144	141	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	145	74	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOTAIL							
126	174	175	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	175	111	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOH							
160	112	110	0	113	0	0	0	0	0	0
0	0.00000	.00000	OO							
33	113	114	0	0	24	0	0	0	0	0
0	0.00000	.00000	OO							
157	114	23	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOLREF							
19	108	109	115	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
134	176	178	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	177	73	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOHEAD							
13	178	177	179	0	0	0	0	0	0	0
0	183.00000	.00000	OO							
7	179	180	181	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	180	33	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOLEADER							
166	181	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	182	74	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOTAIL							
13	183	182	184	0	0	0	0	0	0	0
0	186.00000	.00000	OO							
157	184	73	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOHEAD							
157	185	86	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOZ							
13	186	185	187	0	0	0	0	0	0	0
0	189.00000	.00000	OO							



0	0.00000	.00000	OOY							
157	224	81 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOT							
22	230	231 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
109	238	235 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
95	239	0 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	227	60 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOTRAILER							
5	228	229 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
91	229	230 232	0	0	0	0	0	0	0	0
0	236.00000	.00000	OO							
157	231	64 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOID							
95	232	0 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	235	69 0	0	0	0	0	0	0	0	0
0	239.00000	.00000	OONXT							
62	236	237 238	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
138	237	0 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO=							
157	240	81 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOT							
131	241	240 242	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	242	64 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOID							
13	225	224 226	0	0	0	0	0	0	0	0
0	243.00000	.00000	OO							
7	226	227 228	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
13	243	241 244	0	0	0	0	0	0	0	0
0	248.00000	.00000	OO							
157	244	76 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOQ							
157	245	81 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOT							
131	246	245 247	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	247	69 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OONXT							
13	248	246 250	0	0	0	0	0	0	0	0
0	255.00000	.00000	OO							
157	249	75 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOP							
131	250	249 251	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	251	49 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOTRAIL							
157	252	75 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOP							
131	253	252 254	0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
157	254	49 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOTRAIL							



0	0.00000	.00000	00							
0	134	285	287	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	157	286	76	0	0	0	0	0	0	0
0	0.00000	.00000	00Q							
0	13	287	286	288	0	0	0	0	0	0
0	290.00000	.00000	00							
0	157	288	75	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
0	157	289	75	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
0	13	290	289	292	0	0	0	0	0	0
0	294.00000	.00000	00							
0	157	291	75	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
0	131	292	291	293	0	0	0	0	0	0
0	0.00000	.00000	00							
0	157	293	55	0	0	0	0	0	0	0
0	0.00000	.00000	00NXT							
0	134	303	307	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	157	304	76	0	0	0	0	0	0	0
0	0.00000	.00000	00Q							
0	131	305	304	306	0	0	0	0	0	0
0	0.00000	.00000	00							
0	157	306	55	0	0	0	0	0	0	0
0	0.00000	.00000	00NXT							
0	13	307	305	308	0	0	0	0	0	0
0	310.00000	.00000	00							
0	157	308	73	0	0	0	0	0	0	0
0	0.00000	.00000	00HEAD							
0	157	309	73	0	0	0	0	0	0	0
0	0.00000	.00000	00HEAD							
0	13	310	309	311	0	0	0	0	0	0
0	0.00000	.00000	00							
0	157	311	76	0	0	0	0	0	0	0
0	0.00000	.00000	00Q							
0	71	294	295	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	29	295	299	303	0	0	0	0	0	0
0	0.00000	.00000	00							
0	157	296	76	0	0	0	0	0	0	0
0	0.00000	.00000	00Q							
0	131	297	296	298	0	0	0	0	0	0
0	302.00000	.00000	00							
0	157	298	43	0	0	0	0	0	0	0
0	0.00000	.00000	00COUNT							
0	62	299	300	301	0	0	0	0	0	0
0	0.00000	.00000	00							
0	138	300	0	0	0	0	0	0	0	0
0	0.00000	.00000	00=							
0	109	301	297	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	100	302	4	0	0	0	0	0	0	0
0	0.00000	.00000	100							
0	157	312	76	0	0	0	0	0	0	0
0	0.00000	.00000	00Q							
0	88	278	279	285	0	0	0	0	0	0
0	313.00000	.00000	00							



0	0.00000	.00000	00Q							
131	346	345 347	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	347	55 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00NXT							
134	357	359 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	358	75 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
13	359	358 361	0	0	0	0	0	0	0	0
0	366.00000	.00000	00							
157	360	81 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00T							
131	361	360 362	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	362	64 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00ID							
157	363	75 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
131	364	363 365	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	365	43 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00COUNT							
62	370	371 372	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	371	0 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00-							
109	372	368 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	373	4 0	0	0	0	0	0	0	0	0
0	0.10000+001	.00000	101							
13	366	364 370	0	0	0	0	0	0	0	0
0	374.00000	.00000	00							
157	367	75 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
131	368	367 369	0	0	0	0	0	0	0	0
0	373.00000	.00000	00							
157	369	43 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00COUNT							
134	383	387 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	384	75 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
131	385	384 386	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	386	55 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00NXT							
13	387	385 388	0	0	0	0	0	0	0	0
0	390.00000	.00000	00							
157	388	76 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00Q							
157	389	76 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00Q							
13	390	389 391	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	391	75 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00P							
71	374	375 0	0	0	0	0	0	0	0	0
0	393.00000	.00000	00							



0	0.00000	.00000	00							
0	134 404	406 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	157 405	0 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00PUT							
0	109 407	408 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	119 406	405 407	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	135 408	0 0	0	0	0	0	0	0	0	0
409	0.00000	.00000	00	"THIS SET IS NOT						
0	487 409	0 0	0	0	0	0	0	0	0	0
410	0.00000	.00000	00	PARTIALLY ORDER						
0	487 410	0 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00ED."							
0	157 399	86 0	0	0	0	0	0	0	0	0
0	403.00000	.00000	00Z							
0	62 400	401 402	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	138 401	0 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00=							
0	109 402	399 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	100 403	4 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	100							
0	86 315	317 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00OUTER							
0	71 397	398 0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	29 398	400 404	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	19 16	17 176	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	137 107	97 96	108	0	0	0	0	0	0	0
0	0.00000	.00000	00							
0	27 12	0 0	15	0	0	0	0	0	0	0
0	0.00000	.00000	00							

12.2. APPENDIX IVb.

0	28	11	12	0	0	0	0	0	0	0	0
0	118	0.00000	.00000	00	0	0	0	0	0	0	0
0	117	13	0	0	0	0	0	0	0	0	0
0	117	0.00000	.00000	00	0	0	0	0	0	0	0
0	117	14	13	16	0	0	0	0	0	0	0
0	137	0.00000	.00000	OOINSERT	16	0	0	0	0	0	0
0	137	15	14	13	16	0	0	0	0	0	0
0	82	0.00000	.00000	00	0	0	0	0	0	0	0
0	82	17	18	0	0	0	0	0	0	0	0
0	149	0.00000	.00000	00	0	0	0	0	0	0	0
0	149	18	19	0	0	0	0	0	0	0	0
0	150	20.00000	.00000	00	0	0	0	0	0	0	0
0	150	19	0	0	0	0	0	0	0	0	0
0	149	0.00000	.00000	00NODE	22	0	0	0	0	0	0
0	149	20	21	0	22	0	0	0	0	0	0
0	150	25.00000	.00000	00	0	0	0	0	0	0	0
0	150	21	22	0	0	0	0	0	0	0	0
0	3	0.00000	.00000	00REF	0	0	0	0	0	0	0
0	3	22	23	0	0	0	0	0	0	0	0
0	33	0.00000	.00000	00	0	0	0	0	0	0	0
0	33	23	24	0	0	0	0	0	0	0	0
0	157	0.00000	.00000	00	0	0	0	0	0	0	0
0	157	24	19	0	0	0	0	0	0	0	0
0	123	0.00000	.00000	00NODE	0	0	0	0	0	0	0
0	123	27	28	0	0	0	0	0	0	0	0
0	70	0.00000	.00000	00	0	0	0	0	0	0	0
0	70	29	30	0	0	0	0	0	0	0	0
0	24	0.00000	.00000	00	0	0	0	0	0	0	0
0	24	30	31	0	0	0	0	0	0	0	0
0	33	0.00000	.00000	00KEY	4	4	0	0	0	0	0
0	33	31	32	0	4	4	0	0	0	0	0
0	157	0.00000	.00000	00	0	0	0	0	0	0	0
0	157	32	4	0	0	0	0	0	0	0	0
0	160	0.00000	.00000	00INTEGER	31	0	0	0	0	0	0
0	160	28	29	0	31	0	0	0	0	0	0
0	70	33.00000	.00000	00	0	0	0	0	0	0	0
0	70	34	35	0	0	0	0	0	0	0	0
0	24	0.00000	.00000	00	0	0	0	0	0	0	0
0	24	35	36	0	0	0	0	0	0	0	0
0	33	0.00000	.00000	00COUNT	4	4	0	0	0	0	0
0	33	36	37	0	4	4	0	0	0	0	0
0	157	0.00000	.00000	00	0	0	0	0	0	0	0
0	157	37	4	0	0	0	0	0	0	0	0
0	160	0.00000	.00000	00INTEGER	36	0	0	0	0	0	0
0	160	33	34	0	36	0	0	0	0	0	0
0	70	38.00000	.00000	00	0	0	0	0	0	0	0
0	70	39	40	0	0	0	0	0	0	0	0
0	24	0.00000	.00000	00	0	0	0	0	0	0	0
0	24	40	41	0	0	0	0	0	0	0	0
0	33	0.00000	.00000	00NXT	0	22	0	0	0	0	0
0	33	41	42	0	0	22	0	0	0	0	0
0	157	0.00000	.00000	00	0	0	0	0	0	0	0
0	157	42	21	0	0	0	0	0	0	0	0
0	160	0.00000	.00000	00REF	41	0	0	0	0	0	0
0	160	38	39	0	41	0	0	0	0	0	0
0	0	0.00000	.00000	00	00						

149	25	26	0	27	0	0	0	0	0	0
0	45.00000	.00000	00	00						
150	26	27	0	0	0	0	0	0	0	0
0	0.00000	.00000	00N0D0E	00N0D0E						
70	43	44	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
161	44	46	0	0	0	0	0	0	0	0
0	0.00000	.00000	00K	00K						
160	45	43	0	46	0	0	0	0	0	0
0	50.00000	.00000	00	00						
33	46	47	0	4	4	0	0	0	0	0
0	0.00000	.00000	00	00						
157	47	4	0	0	0	0	0	0	0	0
0	0.00000	.00000	00I0N0T0E0G0E0R	00I0N0T0E0G0E0R						
100	61	4	0	0	0	0	0	0	0	0
0	64.00000	.00000	100	100						
95	64	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
157	54	26	0	0	0	0	0	0	0	0
0	0.00000	.00000	00N0D0E	00N0D0E						
5	55	58	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
100	58	4	0	0	0	0	0	0	0	0
0	61.00000	.00000	100	100						
70	48	49	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
161	49	51	53	0	0	0	0	0	0	0
0	0.00000	.00000	00R000T	00R000T						
110	67	68	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
72	68	69	71	73	0	0	0	0	0	0
0	74.00000	.00000	00	00						
70	69	70	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
73	70	71	73	0	0	0	0	0	0	0
0	0.00000	.00000	00X	00X						
33	71	72	0	4	4	0	0	0	0	0
0	0.00000	.00000	00	00						
157	72	4	0	0	0	0	0	0	0	0
0	0.00000	.00000	00I0N0T0E0G0E0R	00I0N0T0E0G0E0R						
166	73	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
75	74	75	77	79	0	0	0	0	0	0
0	0.00000	.00000	00	00						
70	75	76	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
76	76	77	0	0	0	0	0	0	0	0
0	0.00000	.00000	00B0A0S0E	00B0A0S0E						
33	77	78	0	0	22	0	0	0	0	0
0	0.00000	.00000	00	00						
157	78	21	0	0	0	0	0	0	0	0
0	0.00000	.00000	00R0E0F	00R0E0F						
166	79	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
160	50	48	53	51	0	0	0	0	0	0
0	80.00000	.00000	00	00						
118	65	67	0	0	0	0	0	0	0	0
0	0.00000	.00000	00	00						
117	66	65	81	0	0	0	0	0	0	0

0	0.00000	.00000	OOSEARCH							
0	33 51	52 0	0 22	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	157 52	21 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOREF							
0	7 53	54 55	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	82 82	85 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	70 83	84 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	161 84	86 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOW							
0	160 85	83 0	86 0	0	0	0	0	0	0	0
0	90.00000	.00000	OO							
0	33 86	87 0	0 22	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	157 87	21 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOREF							
0	70 88	89 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	161 89	91 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOB							
0	134 93	95 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	157 94	84 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOW							
0	13 95	94 96	0 0	0	0	0	0	0	0	0
0	98.00000	.00000	OO							
0	157 96	76 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOBASE							
0	157 97	89 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOB							
0	13 98	97 99	0 0	0	0	0	0	0	0	0
0	100.00000	.00000	OO							
0	158 99	0 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	1OTRUE							
0	134 111	112 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	134 121	123 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	157 122	89 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOB							
0	13 123	122 124	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	158 124	0 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	1OFALSE							
0	29 113	117 121	0 0	0	0	0	0	0	0	0
0	125.00000	.00000	OO							
0	29 125	126 127	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	134 127	129 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	157 128	84 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOW							
0	13 129	128 131	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OO							
0	157 130	84 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OOW							



0	0.00000	.00000	00N0D0							
5	144	147	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	147	4	0	0	0	0	0	0	0	0
0	150.00000	.00000	100							
157	154	76	0	0	0	0	0	0	0	0
0	0.00000	.00000	00B0A5E							
131	155	154	156	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	156	30	0	0	0	0	0	0	0	0
0	0.00000	.00000	00K0Y							
13	141	140	142	0	0	0	0	0	0	0
0	157.00000	.00000	00							
7	142	143	144	0	0	0	0	0	0	0
0	0.00000	.00000	00							
13	157	155	158	0	0	0	0	0	0	0
0	162.00000	.00000	00							
157	158	70	0	0	0	0	0	0	0	0
0	0.00000	.00000	00X							
157	159	76	0	0	0	0	0	0	0	0
0	0.00000	.00000	00B0A5E							
131	160	159	161	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	161	35	0	0	0	0	0	0	0	0
0	0.00000	.00000	00C0U0N0T							
13	162	160	163	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	163	4	0	0	0	0	0	0	0	0
0	0.10000+001	.00000	101							
29	134	135	136	0	0	0	0	0	0	0
0	164.00000	.00000	00							
29	164	165	166	0	0	0	0	0	0	0
0	0.00000	.00000	00							
134	166	170	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	167	84	0	0	0	0	0	0	0	0
0	0.00000	.00000	00W							
131	168	167	169	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	169	35	0	0	0	0	0	0	0	0
0	0.00000	.00000	00C0U0N0T							
13	170	168	174	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	171	84	0	0	0	0	0	0	0	0
0	0.00000	.00000	00W							
131	172	171	173	0	0	0	0	0	0	0
0	177.00000	.00000	00							
157	173	35	0	0	0	0	0	0	0	0
0	0.00000	.00000	00C0U0N0T							
62	174	175	176	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	175	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00+							
109	176	172	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	177	4	0	0	0	0	0	0	0	0
0	0.10000+001	.00000	101							
166	165	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							



0	0.00000	.00000		OOCOUNT							
157	210	183	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOV							
13	211	210	213	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	212	183	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOV							
131	213	212	214	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	214	40	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OONXT							
88	190	191	197	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
167	191	193	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	192	183	0	0	0	0	0	0	0	0	0
0	196.00000	.00000		OOV							
62	193	194	195	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
138	194	0	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO/=							
109	195	192	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
95	196	0	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
19	188	0	189	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
134	215	217	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	216	0	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOGET							
109	218	219	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
119	217	216	218	0	0	0	0	0	0	0	0
0	220.00000	.00000		OO							
157	219	44	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOK							
134	227	229	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	228	66	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOSEARCH							
109	230	231	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	231	44	0	0	0	0	0	0	0	0	0
0	232.00000	.00000		OOK							
119	229	228	230	0	0	0	0	0	0	0	0
0	234.00000	.00000		OO							
157	232	49	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOROOT							
157	233	0	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOGET							
109	235	236	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	236	44	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOK							
119	234	233	235	0	0	0	0	0	0	0	0
0	0.00000	.00000		OO							
157	237	179	0	0	0	0	0	0	0	0	0
0	0.00000	.00000		OOPRINTLIST							

	88	220	221	227	0	0	0	0	0	0	0
0	238	.00000		.00000	00						
	167	221	223	0	0	0	0	0	0	0	0
0	0	.00000		.00000	00						
	157	222	44	0	0	0	0	0	0	0	0
0	226	.00000		.00000	00K						
	62	223	224	225	0	0	0	0	0	0	0
0	0	.00000		.00000	00						
	138	224	0	0	0	0	0	0	0	0	0
0	0	.00000		.00000	00/=						
	109	225	222	0	0	0	0	0	0	0	0
0	0	.00000		.00000	00						
	100	226	4	0	0	0	0	0	0	0	0
0	0	.00000		.00000	100						
	109	239	240	0	0	0	0	0	0	0	0
0	0	.00000		.00000	00						
	119	238	237	239	0	0	0	0	0	0	0
0	0	.00000		.00000	00						
	157	240	49	0	0	0	0	0	0	0	0
0	0	.00000		.00000	00ROOT						
	19	16	17	215	0	0	0	0	0	0	0
0	0	.00000		.00000	00						
	137	187	179	178	188	0	0	0	0	0	0
0	0	.00000		.00000	00						
	27	12	0	0	15	0	0	0	0	0	0
0	0	.00000		.00000	00						

12.3. APPENDIX IVc.

0	28	11	12	0	0	0	0	0	0	0	0
0	118	0.00000	0	.00000	00	0	0	0	0	0	0
0	117	13	13	0	00	0	0	0	0	0	0
0	117	14	13	16	00	0	0	0	0	0	0
0	137	0.00000	14	.00000	00V	0	0	0	0	0	0
0	82	15	14	13	16	0	0	0	0	0	0
0	149	0.00000	17	.00000	00	0	0	0	0	0	0
0	149	18	19	0	20	0	0	0	0	0	0
0	150	24.00000	19	.00000	00	0	0	0	0	0	0
0	150	19	20	0	0	0	0	0	0	0	0
0	81	0.00000	20	.00000	00INTEGER	0	0	0	0	0	0
0	122	0.00000	21	.00000	00	0	0	0	0	0	0
0	100	0.00000	22	.00000	00	4	0	0	0	0	0
0	100	22	4	0	0	0	0	0	0	0	0
0	100	0.10000+001	4	.00000	101	0	0	0	0	0	0
0	149	0.50000+002	4	.00000	1050	0	0	0	0	0	0
0	149	24	25	0	26	0	0	0	0	0	0
0	150	33.00000	25	.00000	00	0	0	0	0	0	0
0	49	0.00000	26	.00000	00COLOUR	0	0	0	0	0	0
0	48	0.00000	27	.00000	00	0	0	0	0	0	0
0	48	28.00000	26	.00000	10RED	0	0	0	0	0	0
0	48	28	26	0	0	0	0	0	0	0	0
0	48	29.10000+001	26	.00000	10ORANGE	0	0	0	0	0	0
0	48	29	26	0	0	0	0	0	0	0	0
0	48	30.20000+001	26	.00000	10YELLOW	0	0	0	0	0	0
0	70	0.30000+001	26	.00000	10GREEN	0	0	0	0	0	0
0	161	0.00000	31	.00000	00	0	0	0	0	0	0
0	161	0.00000	32	.00000	00	0	0	0	0	0	0
0	160	0.00000	33	.00000	00ROBOT	34	0	0	0	0	0
0	33	39.00000	31	.00000	00	26	26	26	0	0	0
0	157	0.00000	34	.00000	00	0	0	0	0	0	0
0	70	0.00000	35	.00000	00COLOUR	0	0	0	0	0	0
0	161	0.00000	36	.00000	00	0	0	0	0	0	0
0	161	38.00000	37	.00000	00VAR2	0	0	0	0	0	0
0	161	0.00000	40	.00000	00VAR1	0	0	0	0	0	0
0	134	0.00000	42	.00000	00	0	0	0	0	0	0
0	87	0.00000	43	.00000	00	0	0	0	0	0	0
0	87	43	44	46	0	0	0	0	0	0	0
0	48	48.00000	44	.00000	00	0	0	0	0	0	0

0	86	44	46	0	0	0	0	0	0	0
0	0.00000		.00000	00LAB1						
0	157	45	38	0	0	0	0	0	0	0
0	0.00000		.00000	00VAR1						
0	13	46	45	47	0	0	0	0	0	0
0	0.00000		.00000	00						
0	100	47	4	0	0	0	0	0	0	0
0	0.20000+001		.00000	102						
0	87	48	49	52	0	0	0	0	0	0
0	55.00000		.00000	00						
0	86	49	52	0	0	0	0	0	0	0
0	50.00000		.00000	00LAB2						
0	86	50	52	0	0	0	0	0	0	0
0	0.00000		.00000	00LAB3						
0	157	51	37	0	0	0	0	0	0	0
0	0.00000		.00000	00VAR2						
0	13	52	51	53	0	0	0	0	0	0
0	0.00000		.00000	00						
0	157	53	38	0	0	0	0	0	0	0
0	0.00000		.00000	00VAR1						
0	157	54	37	0	0	0	0	0	0	0
0	0.00000		.00000	00VAR2						
0	13	55	54	56	0	0	0	0	0	0
0	57.00000		.00000	00						
0	100	56	4	0	0	0	0	0	0	0
0	0.10000+002		.00000	1010						
0	97	57	0	0	0	0	0	0	0	0
0	58.00000		.00000	00						
0	69	58	59	0	0	0	0	0	0	0
0	60.00000		.00000	00						
0	157	59	49	0	0	0	0	0	0	0
0	0.00000		.00000	00LAB2						
0	157	61	44	0	0	0	0	0	0	0
0	0.00000		.00000	00LAB1						
0	69	60	61	0	0	0	0	0	0	0
0	62.00000		.00000	00						
0	69	62	63	0	0	0	0	0	0	0
0	64.00000		.00000	00						
0	157	63	50	0	0	0	0	0	0	0
0	0.00000		.00000	00LAB3						
0	134	70	71	0	0	0	0	0	0	0
0	0.00000		.00000	00						
0	9	72	74	0	0	0	0	0	0	0
0	0.00000		.00000	00						
0	157	73	38	0	0	0	0	0	0	0
0	0.00000		.00000	00VAR1						
0	22	75	76	0	0	0	0	0	0	0
0	0.00000		.00000	00						
0	134	78	79	0	0	0	0	0	0	0
0	0.00000		.00000	00						
0	82	80	83	0	0	0	0	0	0	0
0	0.00000		.00000	00						
0	70	81	82	0	0	0	0	0	0	0
0	0.00000		.00000	00						
0	161	82	84	0	0	0	0	0	0	0
0	0.00000		.00000	00VAR4						
0	134	86	88	0	0	0	0	0	0	0
0	0.00000		.00000	00						
0	157	87	82	0	0	0	0	0	0	0

0	0.00000	.00000	OOVAR4							
13	88	87	90	0	0	0	0	0	0	0
0	95.00000	.00000	00							
100	89	3	0	0	0	0	0	0	0	0
0	93.45000+001.	.00000	104.5							
62	90	91	92	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	91	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00*							
109	92	89	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	93	3	0	0	0	0	0	0	0	0
0	0.33000+001.	.00000	103.3							
157	94	82	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOVAR4							
13	95	94	97	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	96	82	0	0	0	0	0	0	0	0
0	100.00000	.00000	OOVAR4							
62	97	98	99	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	98	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00**							
109	99	96	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	100	4	0	0	0	0	0	0	0	0
0	0.20000+001.	.00000	102							
19	79	80	86	0	0	0	0	0	0	0
0	0.00000	.00000	00							
160	83	81	0	84	0	0	0	0	0	0
0	0.00000	.00000	00							
33	84	85	0	3	3	0	0	0	0	0
0	0.00000	.00000	00							
157	85	3	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOFLOAT							
8	74	75	78	0	0	0	0	0	0	0
0	101.00000	.00000	00							
100	76	4	0	0	0	0	0	0	0	0
0	0.10000+001.	.00000	101							
100	77	4	0	0	0	0	0	0	0	0
0	0.20000+001.	.00000	102							
22	102	103	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
134	104	105	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
134	107	109	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	108	38	0	0	0	0	0	0	0	0
0	0.00000	.00000	OOVAR1							
13	109	108	111	0	0	0	0	0	0	0
0	115.00000	.00000	00							
157	110	38	0	0	0	0	0	0	0	0
0	114.00000	.00000	OOVAR1							
62	111	112	113	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	112	0	0	0	0	0	0	0	0	0
0	0.00000	.00000	00+							
109	113	110	0	0	0	0	0	0	0	0
0	0.00000	.00000	00							



0	0.00000	.00000	OONAMED1							
92	67	66 68	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
88	68	69 70	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
166	69	0 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	144	38 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00VAR1							
134	143	145 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
13	145	144 147	0 0	0	0	0	0	0	0	0
0	152.00000	.00000	00							
157	146	38 0	0 0	0	0	0	0	0	0	0
0	150.00000	.00000	00VAR1							
62	147	148 149	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	148	0 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00+							
109	149	146 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	150	4 0	0 0	0	0	0	0	0	0	0
0	0.10000+001	.00000	101							
157	151	37 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00VAR2							
13	152	151 154	0 0	0	0	0	0	0	0	0
0	158.00000	.00000	00							
157	153	37 0	0 0	0	0	0	0	0	0	0
0	157.00000	.00000	00VAR2							
62	154	155 156	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	155	0 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00+							
109	156	153 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	157	38 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00VAR1							
52	158	159 161	138 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
157	159	136 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	OONAMED2							
157	160	38 0	0 0	0	0	0	0	0	0	0
0	164.00000	.00000	00VAR1							
62	161	162 163	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
138	162	0 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00=							
109	163	160 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
100	164	4 0	0 0	0	0	0	0	0	0	0
0	0.30000+001	.00000	103							
92	137	136 138	0 0	0	0	0	0	0	0	0
0	166.00000	.00000	00							
88	138	139 143	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
56	139	140 141	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00							
83	140	141 0	0 0	0	0	0	0	0	0	0
0	0.00000	.00000	00I							

0	33	141	142	0	26	26	26	0	0	0	0
0	157	142	25	0	00	0	0	0	0	0	0
0	82	168	171	0	00	0	0	0	0	0	0
0	70	169	170	0	00	0	0	0	0	0	0
0	161	170	172	0	00	0	0	0	0	0	0
0	134	174	176	0	00	0	0	0	0	0	0
0	157	175	170	0	00	0	0	0	0	0	0
0	13	176	175	177	00	0	0	0	0	0	0
0	100	177	3	0	00	0	0	0	0	0	0
0	134	185	187	0	103.3	0	0	0	0	0	0
0	157	186	170	0	00	0	0	0	0	0	0
0	100	188	3	0	00	0	0	0	0	0	0
0	13	187	186	188	106.6	0	0	0	0	0	0
0	71	178	179	0	00	0	0	0	0	0	0
0	29	179	181	185	00	0	0	0	0	0	0
0	157	180	38	0	00	0	0	0	0	0	0
0	62	181	182	183	00	0	0	0	0	0	0
0	138	182	0	0	00	0	0	0	0	0	0
0	109	183	180	0	00=	0	0	0	0	0	0
0	157	184	37	0	00	0	0	0	0	0	0
0	92	166	165	167	00	0	0	0	0	0	0
0	19	167	168	174	00	0	0	0	0	0	0
0	160	171	169	0	00	172	0	0	0	0	0
0	33	172	173	0	00	3	3	0	0	0	0
0	157	173	3	0	00	0	0	0	0	0	0
0	134	196	197	0	00	0	0	0	0	0	0
0	134	204	206	0	00	0	0	0	0	0	0
0	157	205	38	0	00	0	0	0	0	0	0
0	13	206	205	207	00	0	0	0	0	0	0
0	100	207	4	0	00	0	0	0	0	0	0





0	0.00000	.00000	00ROBOT							
86	165	167	0	0	0	0	0	0	0	0
0	0.00000	.00000	00NAMED3							
13	253	252	254	0	0	0	0	0	0	0
0	0.00000	.00000	00							
86	136	138	0	0	0	0	0	0	0	0
0	0.00000	.00000	00NAMED2							
158	254	30	26	0	0	0	0	0	0	0
0	0.00000	.00000	10GREEN							
19	16	17	42	0	0	0	0	0	0	0
0	0.00000	.00000	00							
160	39	36	0	40	0	0	0	0	0	0
0	0.00000	.00000	00							
33	40	41	0	4	20	21	0	0	0	0
0	0.00000	.00000	00							
157	41	19	0	0	0	0	0	0	0	0
0	0.00000	.00000	00INTEGER							
27	12	0	0	15	0	0	0	0	0	0
0	0.00000	.00000	00							

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